

Combinatorics

Extremal Graph Theory

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Extremal Combinatorics

“how large or how small a collection of finite objects can be, if it has to satisfy certain restrictions”

Extremal Problem:

“What is the largest number of edges that an n -vertex *cycle-free* graph can have?”

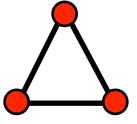
$$(n - 1)$$

Extremal Graph:

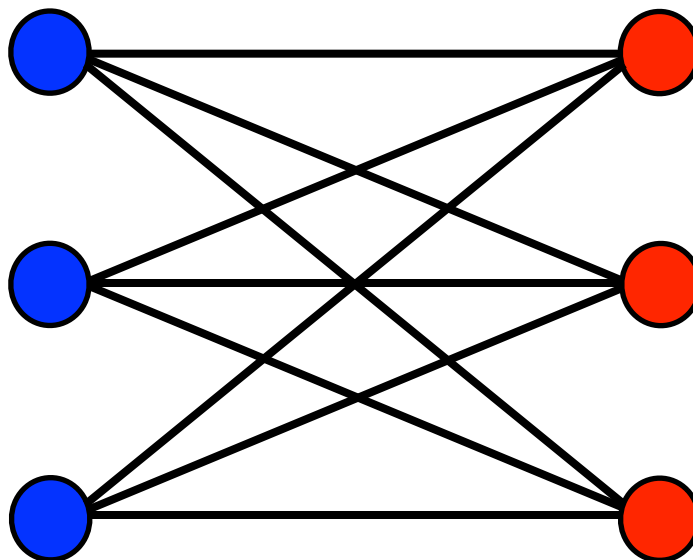
spanning tree

Triangle-Freeness

Triangle-free graph

contains no  as subgraph

Example: bipartite graph



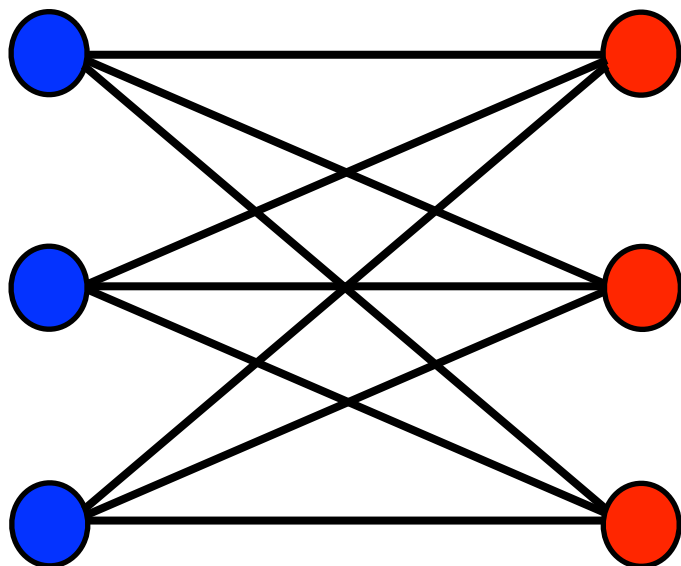
$|E|$ is maximized for
complete balanced bipartite graph

Extremal?

Mantel's Theorem

Theorem (Mantel 1907)

If $G(V, E)$ has $|V| = n$ and is **triangle-free**,
then $|E| \leq \frac{n^2}{4}$.



For n is even,
extremal graph:

$$K_{\frac{n}{2}, \frac{n}{2}}$$

$$\triangle\text{-free} \implies |E| \leq n^2/4$$

First Proof. Induction on n .

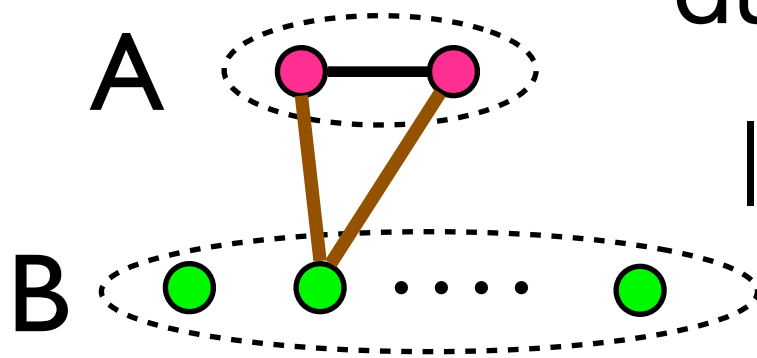
Basis: $n = 1, 2$. trivial

Induction Hypothesis: for any $n < N$

$$|E| > \frac{n^2}{4} \implies G \supseteq \triangle$$

Induction step: for $n = N$

due to **I.H.** $|E(B)| \leq (n-2)^2/4$



$$|E(A, B)| = |E| - |E(B)| - 1$$

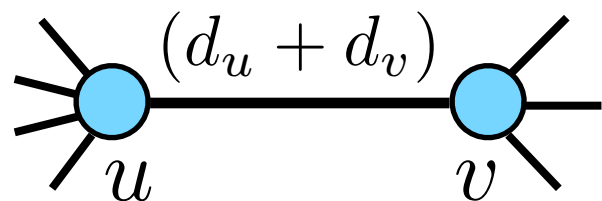
$$> \frac{n^2}{4} - \frac{(n-2)^2}{4} - 1 = n - 2$$

pigeonhole!

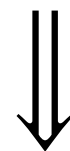
$$\triangle\text{-free} \implies |E| \leq n^2/4$$

Second Proof.

\triangle -free



$$\implies d_u + d_v \leq n, \quad \forall uv \in E$$



Double counting: $\sum_{v \in V} d_v^2 = \sum_{uv \in E} (d_u + d_v) \leq n |E|$

Cauchy-Schwarz

(handshaking)

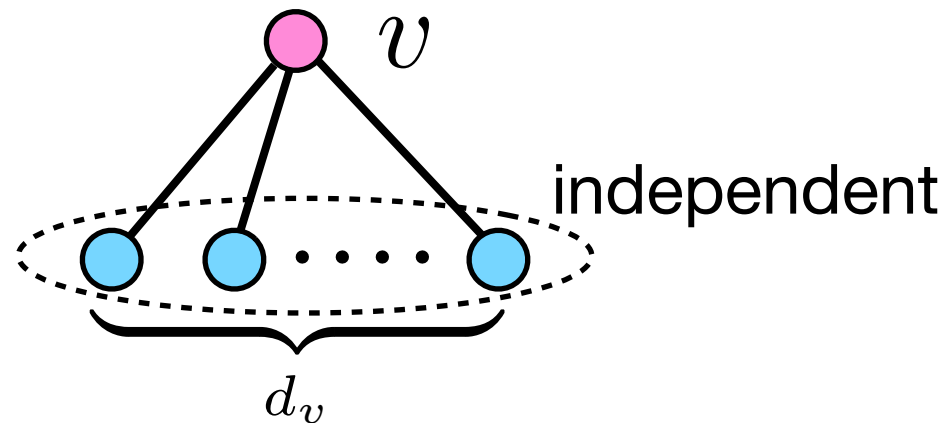
$$n^2 |E| \geq n \sum_{v \in V} d_v^2 = \left(\sum_{v \in V} 1^2 \right) \left(\sum_{v \in V} d_v^2 \right) \geq \left(\sum_{v \in V} d_v \right)^2 = 4 |E|^2$$

$$\implies |E| \leq n^2/4$$

$$\triangle\text{-free} \implies |E| \leq n^2/4$$

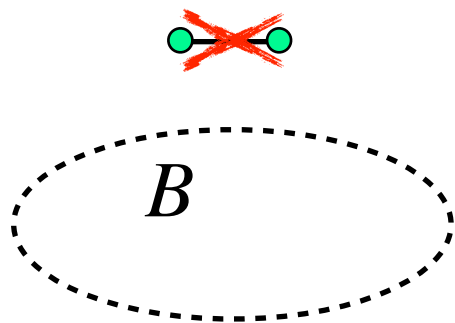
Third Proof.

A : maximum independent set $\alpha = |A|$



$$\forall v \in V, d_v \leq \alpha$$

$B = V \setminus A$ B incident to all edges $\beta = |B|$



Inequality of the arithmetic and geometric mean

$$|E| \leq \sum_{v \in B} d_v \leq \alpha \beta \leq \left(\frac{\alpha + \beta}{2} \right)^2 = \frac{n^2}{4}$$

Turán's Theorem



Paul Turán
(1910-1976)

Turán's Theorem

*“Suppose G is a K_r -free graph.
What is the largest number of
edges that G can have?”*

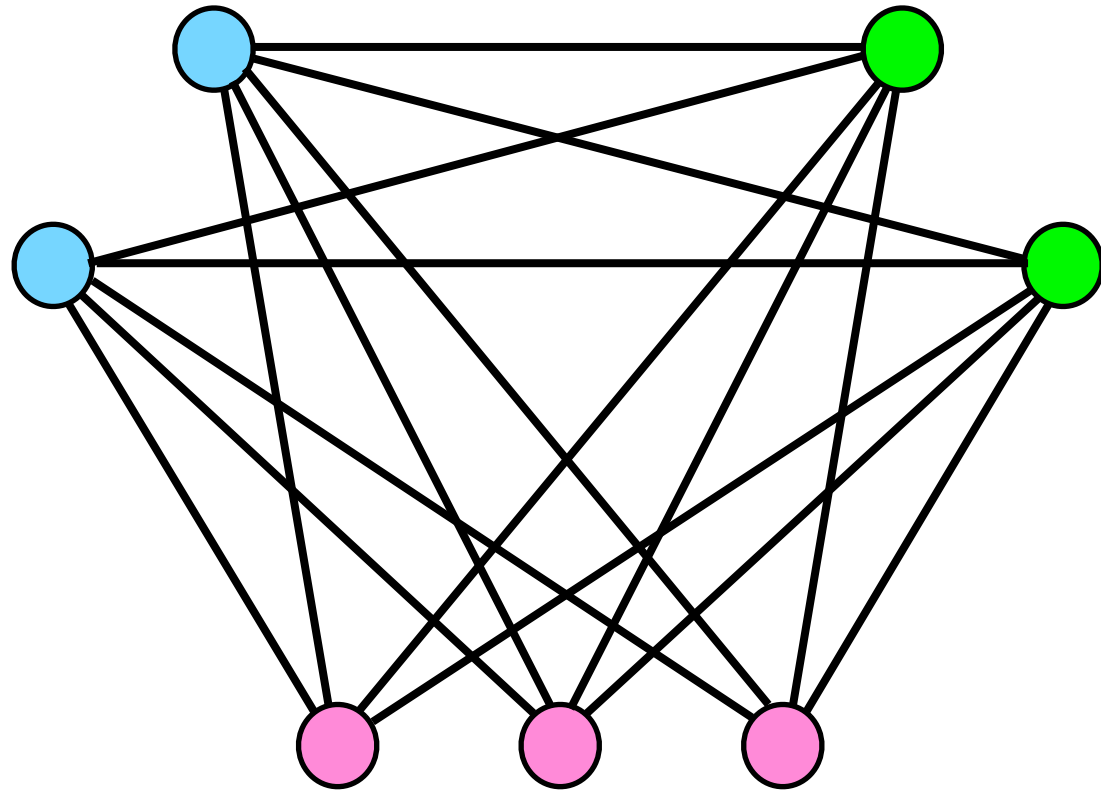
Theorem (Turán 1941)

If $G(V, E)$ has $|V| = n$ and is K_r -free, then

$$|E| \leq \frac{r-2}{2(r-1)} n^2$$

Complete multipartite graph K_{n_1, n_2, \dots, n_r}

$K_{2,2,3}$



Turán graph $T(n, r)$:

$$T(n, r) = K_{n_1, n_2, \dots, n_r}$$

where $n_1 + \dots + n_r = n$ and $n_i \in \left\{ \left\lfloor \frac{n}{r} \right\rfloor, \left\lceil \frac{n}{r} \right\rceil \right\}$

Turán graph $T(n, r)$:

$$T(n, r) = K_{n_1, n_2, \dots, n_r}$$

where $n_1 + \dots + n_r = n$ and $n_i \in \left\{ \left\lfloor \frac{n}{r} \right\rfloor, \left\lceil \frac{n}{r} \right\rceil \right\}$

$T(n, r - 1)$ has no K_r

$$\begin{aligned} |T(n, r - 1)| &\leq \binom{r - 1}{2} \left(\frac{n}{r - 1} \right)^2 \\ &= \frac{r - 2}{2(r - 1)} n^2 \end{aligned}$$

$$K_r\text{-free} \implies |E| \leq \frac{r-2}{2(r-1)} n^2$$

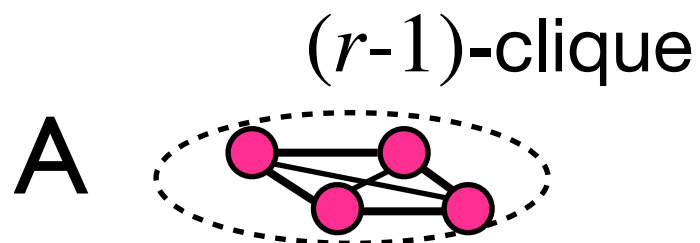
First Proof. (Induction)

Basis: $n = 1, 2, \dots, r-1$.

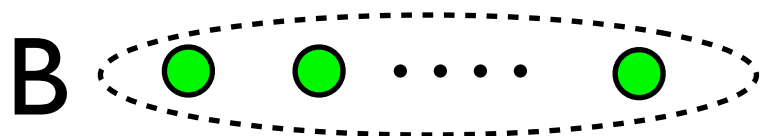
Induction Hypothesis: true for any $n < N$

Induction step: for $n = N$,

suppose G is maximum K_r -free



$\exists (r-1)$ -clique

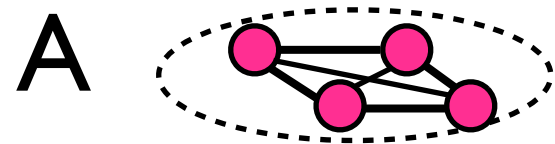


$$K_r\text{-free} \implies |E| \leq \frac{r-2}{2(r-1)}n^2$$

First Proof. (Induction)

suppose G is maximum K_r -free

($r-1$)-clique **I.H.:** $|E(B)| \leq \frac{r-2}{2(r-1)}(n-r+1)^2$



K_r -free \implies no $u \in B \sim$ all $v \in A$

B  $\implies |E(A, B)| \leq (r-2)(n-r+1)$

$$|E| = |E(A)| + |E(B)| + |E(A, B)|$$

$$\leq \binom{r-1}{2} + \frac{r-2}{2(r-1)}(n-r+1)^2 + (r-2)(n-r+1)$$

$$= \frac{r-2}{2(r-1)}n^2$$

$$K_r\text{-free} \implies |E| \leq \frac{r-2}{2(r-1)} n^2$$

Second Proof. (weight shifting)

Assign each vertex v a **weight** $w_v > 0$ s.t. $\sum_{v \in V} w_v = 1$

Evaluate $S(\vec{w}) = \sum_{uv \in E} w_u w_v$

Let $W_u = \sum_{v \sim u} w_v$ For $u \sim v$ that $W_u \geq W_v$

$$(w_u + \epsilon)W_u + (w_v - \epsilon) \geq w_u W_u + w_v W_v$$

shifting all weight of v to $u \implies S(\vec{w})$ non-decreasing

$S(\vec{w})$ is maximized \implies all weights on a clique

$$K_r\text{-free} \implies |E| \leq \frac{r-2}{2(r-1)} n^2$$

Second Proof. (weight shifting)

Assign each vertex v a **weight** $w_v > 0$ s.t. $\sum_{v \in V} w_v = 1$

$$\text{Evaluate } S(\vec{w}) = \sum_{uv \in E} w_u w_v \leq \binom{r-1}{2} \frac{1}{(r-1)^2}$$

$S(\vec{w})$ is maximized \implies all weights on a clique

$$\text{when all } w_v = \frac{1}{n}$$

$$S(\vec{w}) = \sum_{uv \in E} w_u w_v = \frac{|E|}{n^2}$$

$$K_r\text{-free} \implies |E| \leq \frac{r-2}{2(r-1)} n^2$$

Third Proof. (The probabilistic method)

clique number $\omega(G)$: size of the largest clique

$$\omega(G) \geq \sum_{v \in V} \frac{1}{n - d_v}$$

random permutation π of V

$S = \{v \mid \pi_u < \pi_v \implies u \sim v\}$
is a clique

Linearity of expectation:

$$\begin{aligned} \mathbb{E}[|S|] &= \sum_{v \in V} \Pr[v \in S] \geq \sum_{v \in V} \Pr[\forall u \not\sim v : \pi_u \geq \pi_v] \\ &= \sum_{v \in V} \frac{1}{n - d_v} \end{aligned}$$

$$K_r\text{-free} \implies |E| \leq \frac{r-2}{2(r-1)} n^2$$

Third Proof. (The probabilistic method)

$$\omega(G) \geq \sum_{v \in V} \frac{1}{n - d_v}$$

Cauchy-Schwarz

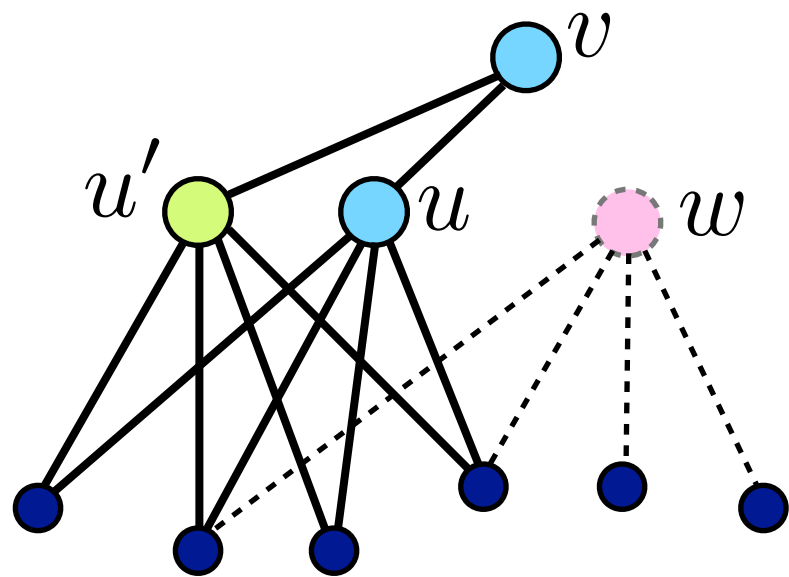
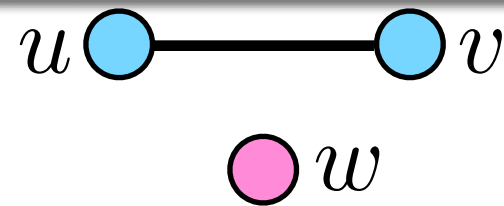
$$\begin{aligned} n = \sum_{v \in V} 1 &\leq \left(\sum_{v \in V} \frac{1}{n - d_v} \right) \left(\sum_{v \in V} (n - d_v) \right) \\ &\leq \omega(G) \sum_{v \in V} (n - d_v) = (r-1)(n^2 - 2|E|) \\ &\quad \text{(handshaking)} \\ &\implies |E| \leq \frac{r-2}{2(r-1)} n^2 \end{aligned}$$

$$K_r\text{-free} \implies |E| \leq \frac{r-2}{2(r-1)} n^2$$

Fourth Proof.

Suppose G is K_r -free with **maximum** edges.

G does not have



By contradiction.

Case.1 $d_w < d_u$ or $d_w < d_v$

duplicate u , delete w , **still K_r -free**

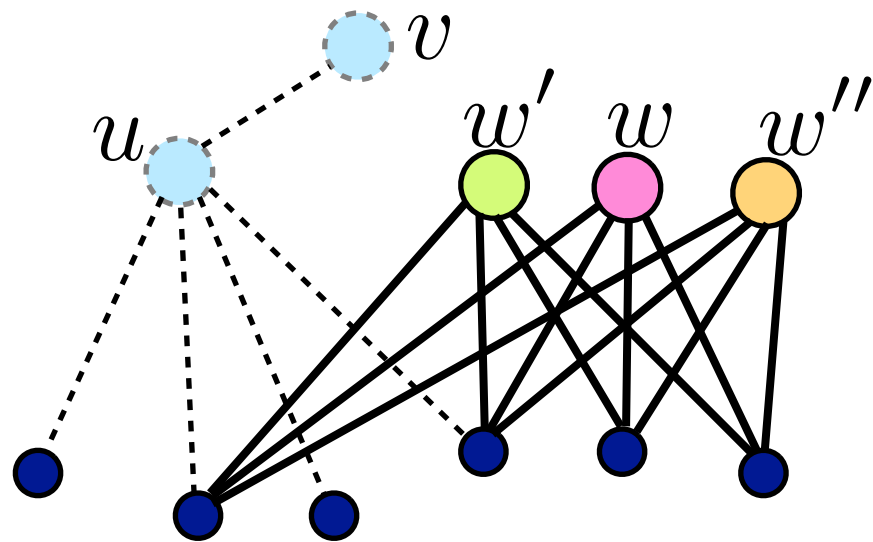
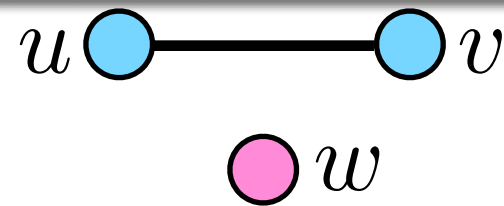
$$|E'| = |E| + d_u - d_w > |E|$$

$$K_r\text{-free} \implies |E| \leq \frac{r-2}{2(r-1)} n^2$$

Fourth Proof.

Suppose G is K_r -free with **maximum** edges.

G does not have



Case.2 $d_w \geq d_u \wedge d_w \geq d_v$

delete u, v , duplicate w , twice

still K_r -free

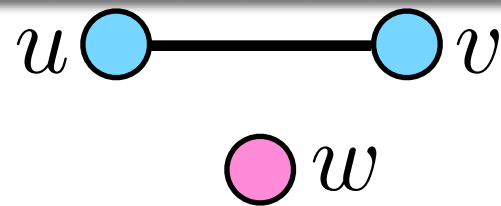
$$|E'| = |E| + 2d_w - (d_u + d_v - 1) > |E|$$

$$K_r\text{-free} \implies |E| \leq \frac{r-2}{2(r-1)} n^2$$

Fourth Proof.

Suppose G is K_r -free with **maximum** edges.

G does not have



$u \approx v$ is an equivalence relation

G is a complete multipartite graph

optimize $K_{n_1, n_2, \dots, n_{r-1}}$

subject to $n_1 + n_2 + \dots + n_r = n$

Turán's Theorem (clique)

If $G(V, E)$ has $|V| = n$ and is K_r -free, then

$$|E| \leq \frac{r-2}{2(r-1)} n^2$$

Turán's Theorem (independent set)

If $G(V, E)$ has $|V| = n$ and $|E| = m$, then
 G has an independent set of size

$$\geq \frac{n^2}{2m + n}$$

Parallel Max

- compute max of n distinct numbers
 - computation model: **parallel, comparison-based**
- 1-round algorithm: $\binom{n}{2}$ comparisons of all pairs
- lower bound for one-round:
 - $\binom{n}{2}$ comparisons are required in the worst case



adversary argument

Parallel Max

- 2-round algorithm:
 - divide n numbers into k groups of n/k each
 - *1st round*: find max of each group;
 $k \binom{n/k}{2}$ comparisons
 - *2nd round*: find the max of the k maxes
 $\binom{k}{2}$ comparisons
- total comparisons: $k \binom{n/k}{2} + \binom{k}{2} = O(n^{4/3})$
for $k = n^{2/3}$

3-round?

optimal?

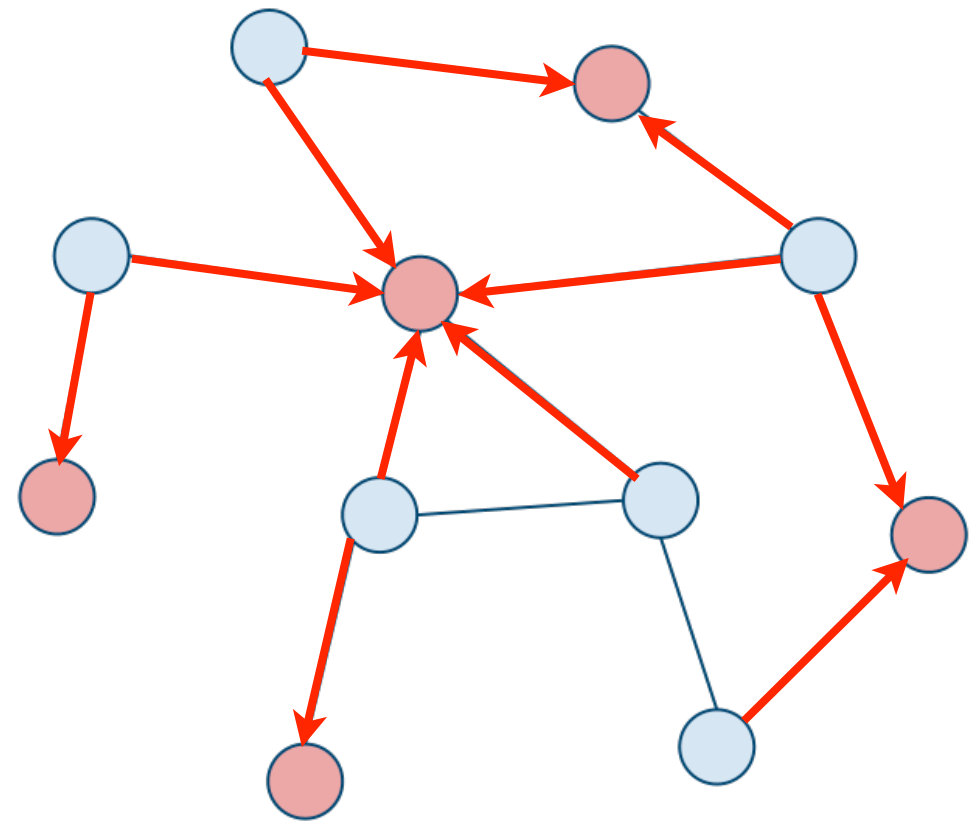
1st round:

Alg: m comparisons



choose an independent set
of size $\geq \frac{n^2}{2m+n}$ (Turán)

make them local maximal



2nd round:

a parallel max problem of size $\geq \frac{n^2}{2m+n}$

requires $\geq \binom{\frac{n^2}{2m+n}}{2}$ comparisons

total comparisons $\geq m + \binom{\frac{n^2}{2m+n}}{2} = \Omega(n^{4/3})$

Fundamental Theorem **of Extremal Graph Theory**

Extremal Graph Theory

Fix a graph H .

$$\text{ex}(n, H)$$

largest possible number of edges
of $G \not\supseteq H$ on n vertices

$$\text{ex}(n, H) = \max_{\substack{G \not\supseteq H \\ |V(G)|=n}} |E(G)|$$

Turán's Theorem

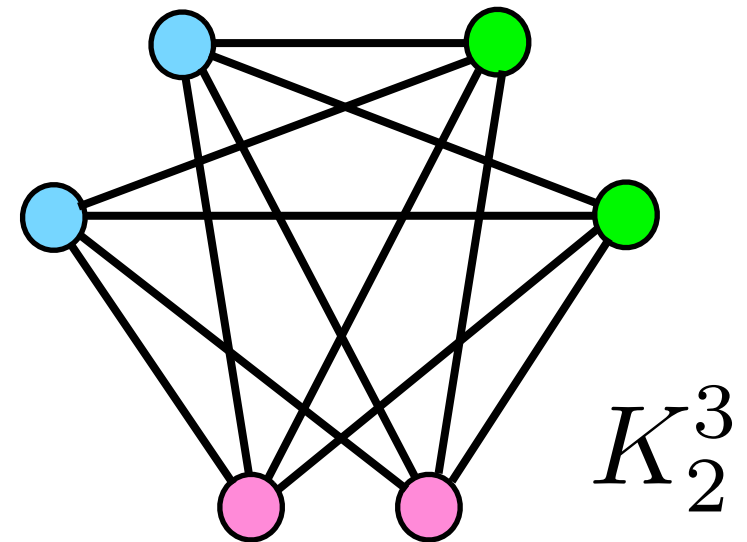
$$\text{ex}(n, K_r) = |T(n, r-1)| \leq \frac{r-2}{2(r-1)} n^2$$

Erdős–Stone theorem

(Fundamental theorem of extremal graph theory)

$$K_s^r = K_{\underbrace{s, s, \dots, s}_r} = T(rs, r)$$

complete r -partite graph
with s vertices in each part



Theorem (Erdős–Stone 1946)

$$\text{ex}(n, K_s^r) = \left(\frac{r-2}{2(r-1)} + o(1) \right) n^2$$

Theorem (Erdős–Stone 1946)

$$\text{ex}(n, K_s^r) = \left(\frac{r-2}{2(r-1)} + o(1) \right) n^2$$

$\text{ex}(n, H) / \binom{n}{2}$ **extremal density** of subgraph H

Corollary

For any nonempty graph H

$$\lim_{n \rightarrow \infty} \frac{\text{ex}(n, H)}{\binom{n}{2}} = \frac{\chi(H) - 2}{\chi(H) - 1}$$

$$\lim_{n \rightarrow \infty} \frac{\text{ex}(n, H)}{\binom{n}{2}} = \frac{\chi(H) - 2}{\chi(H) - 1}$$

$$\chi(H) = r$$

$$H \not\subseteq T(n, r - 1) \text{ for any } n$$

$$\text{ex}(n, H) \geq |T(n, r - 1)|$$

$$H \subseteq K_s^r \text{ for sufficiently large } s$$

$$\text{ex}(n, H) \leq \text{ex}(n, K_s^r)$$

$$= \left(\frac{r-2}{2(r-1)} + o(1) \right) n^2$$

$$\lim_{n \rightarrow \infty} \frac{\text{ex}(n, H)}{\binom{n}{2}} = \frac{\chi(H) - 2}{\chi(H) - 1}$$

$$\chi(H) = r$$

$$|T(n, r - 1)| \leq \text{ex}(n, H) \leq \left(\frac{r - 2}{2(r - 1)} + o(1) \right) n^2$$

$$\frac{r - 2}{r - 1} - o(1) \leq \frac{\text{ex}(n, H)}{\binom{n}{2}} \leq \frac{r - 2}{r - 1} + o(1)$$

Cycles

Girth

girth $g(G)$: length of the shortest cycle in G

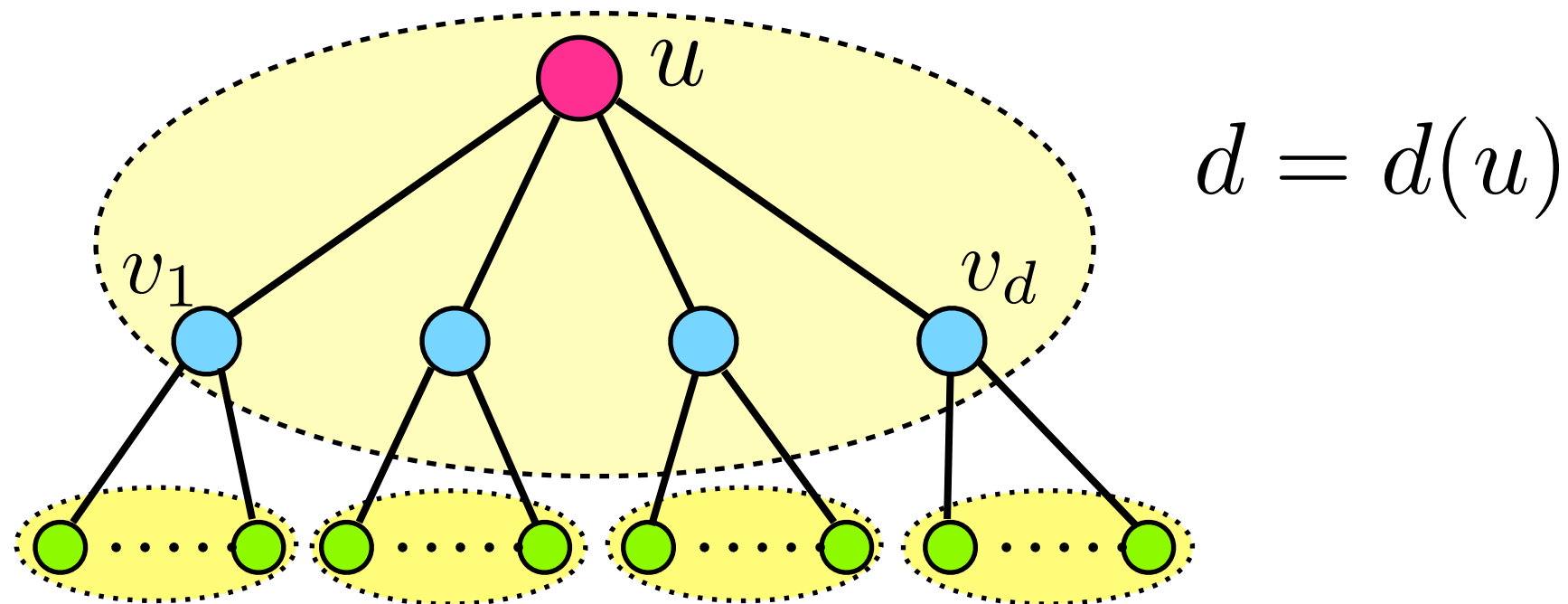
Theorem

If $G(V, E)$ has $|V| = n$ and $\text{girth } g(G) \geq 5$,

$$\text{then } |E| \leq \frac{1}{2}n\sqrt{n-1}$$

$g(G) \geq 5 \iff \triangle\text{- and } \square\text{-free}$

$$g(G) \geq 5 \Rightarrow |E| \leq \frac{1}{2}n\sqrt{n-1}$$



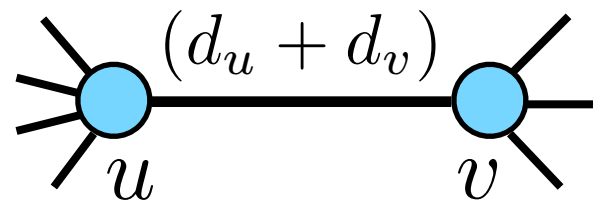
disjoint sets

$$(d+1) + (d(v_1) - 1) + \cdots + (d(v_d) - 1) \leq n$$

$$\sum_{v: v \sim u} d(v) \leq n - 1$$

$$g(G) \geq 5 \Rightarrow |E| \leq \frac{1}{2}n\sqrt{n-1}$$

$$\forall u \in V, \sum_{v:v \sim u} d(v) \leq n-1$$



$$n(n-1) \geq \sum_{u \in V} \sum_{v:v \sim u} d(v) = \sum_{v \in V} d(v)^2$$

$$\geq \frac{\left(\sum_{v \in V} d(v)\right)^2}{n} = \frac{4|E|^2}{n}$$

Cauchy-Schwarz

Hamiltonian Cycle

Dirac's Theorem

$$\forall v \in V, d_v \geq \frac{n}{2} \Rightarrow G(V, E) \text{ is Hamiltonian.}$$

By contradiction, suppose G is the **maximum non-Hamiltonian** graph with $\forall v \in V, d_v \geq \frac{n}{2}$

adding 1 edge \implies Hamiltonian

\exists a Hamiltonian path

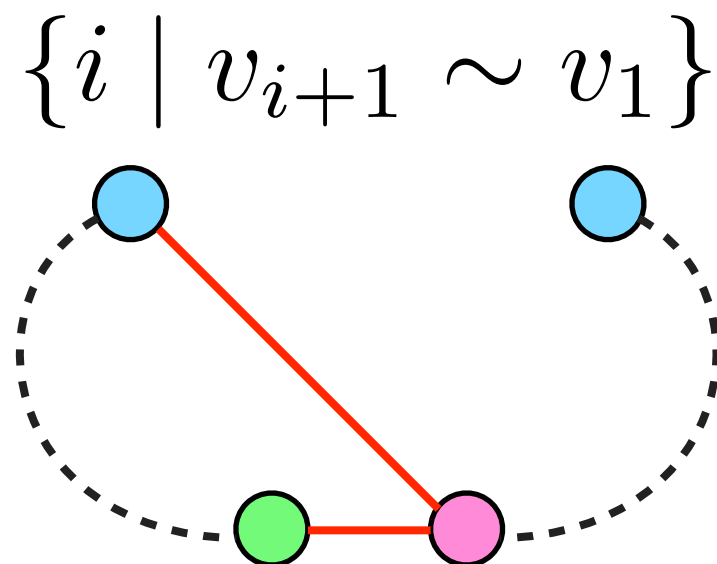
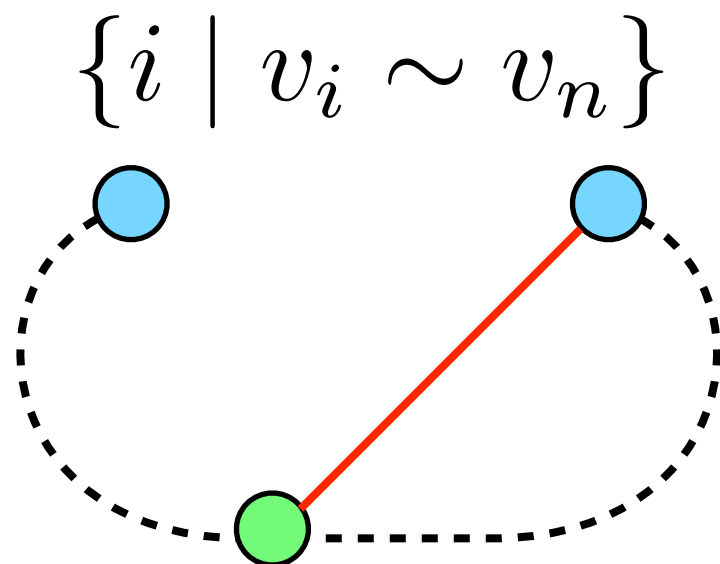
say $v_1 v_2 \cdots v_n$

G is non-Hamiltonian

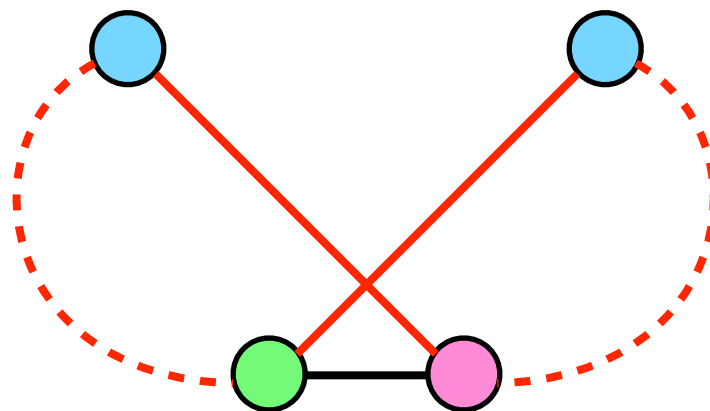
$$\forall v \in V, d_v \geq \frac{n}{2}$$

\exists a Hamiltonian path

$$v_1 v_2 \cdots v_n$$



$\geq \frac{n}{2} + \frac{n}{2}$ pigeons in $\{1, 2, \dots, n-1\}$



Contradiction!