

Randomized Algorithms

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Martingales

Definition:

A sequence of random variables $X_0, X_1, \dots$ is a **martingale** if for all $i > 0$,

$$\mathbf{E}[X_i \mid X_0, \dots, X_{i-1}] = X_{i-1}$$

$$\forall x_0, x_1, \dots, x_{i-1},$$

$$\mathbf{E}[X_i \mid X_0 = x_0, X_1 = x_1, \dots, X_{i-1} = x_{i-1}] = x_{i-1}$$

Azuma's Inequality:

Let $X_0, X_1, \dots$ be a martingale such that, for all $k \geq 1$,

$$|X_k - X_{k-1}| \leq c_k,$$

Then

$$\Pr [|X_n - X_0| \geq t] \leq 2 \exp \left(- \frac{t^2}{2 \sum_{k=1}^n c_k^2} \right).$$

- For a sequence, if in each step:
 - averagely no change to the current value (martingale),
 - no big jump,
- the final does not deviate far from the initial.

Azuma's Inequality:

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1. Represent total difference as sum of step-wise differences.

$$\text{Let } Y_i = X_i - X_{i-1}. \quad X_n - X_0 = \sum_{i=1}^n Y_i$$

2. Apply Markov's inequality to the moment generating function.

$$\Pr[\sum_{i=1}^n Y_i \geq t] = \Pr[e^{\lambda \sum_{i=1}^n Y_i} \geq e^{\lambda t}] \leq \frac{\mathbb{E}[e^{\lambda \sum_{i=1}^n Y_i}]}{e^{\lambda t}}$$

3. Bound the moment generating function.

by **martingale property** & **convexity of MGF**

Generalization

Definition:

$Y_0, Y_1, \dots$ is a martingale with respect to $X_0, X_1, \dots$ if, for all $i \geq 0$,

- Y_i is a function of $X_0, X_1, \dots, X_i$;
- $\mathbf{E}[Y_{i+1} \mid X_0, \dots, X_i] = Y_i$.

Azuma's Inequality (general version):

Let $Y_0, Y_1, \dots$ be a martingale with respect to $X_0, X_1, \dots$ such that, for all $k \geq 1$,

$$|Y_k - Y_{k-1}| \leq c_k,$$

Then

$$\Pr [|Y_n - Y_0| \geq t] \leq 2 \exp \left(- \frac{t^2}{2 \sum_{k=1}^n c_k^2} \right).$$

Doob Sequence

Definition (Doob sequence):

The Doob sequence of a function f with respect to a sequence $X_1, \dots, X_n$ is

$$Y_i = \mathbf{E}[f(X_1, \dots, X_n) \mid X_1, \dots, X_i]$$

$$Y_0 = \mathbf{E}[f(X_1, \dots, X_n)] \dashrightarrow Y_n = f(X_1, \dots, X_n)$$

Doob sequence:

$$Y_i = \mathbf{E}[f(X_1, \dots, X_n) \mid X_1, \dots, X_i]$$

Doob sequence is a martingale:

$$\mathbf{E}[Y_i \mid X_1, \dots, X_{i-1}] = Y_{i-1}$$

Proof:

$$\begin{aligned} & \mathbf{E}[Y_i \mid X_1, \dots, X_{i-1}] \\ &= \mathbf{E}[\mathbf{E}[f(X_1, \dots, X_n) \mid X_1, \dots, X_i] \mid X_1, \dots, X_{i-1}] \\ &= \mathbf{E}[f(X_1, \dots, X_n) \mid X_1, \dots, X_{i-1}] \\ &= Y_{i-1} \end{aligned}$$

Doob Sequence

randomized by

$$f(\textcircled{1}, \underbrace{\textcircled{\$}, \textcircled{\$}, \textcircled{\$}, \textcircled{\$}, \textcircled{\$}}_{\text{averaged over}})$$

Doob Sequence

randomized by

$$f(\overbrace{1, 0}, \underbrace{\$, \$, \$, \$}_{\text{averaged over}})$$

averaged over

Doob Sequence

randomized by

$$f(\underbrace{(1, 0, 0)}_{\text{randomized by}}, \underbrace{(\text{coin}, \text{coin}, \text{coin})}_{\text{averaged over}})$$

averaged over

Doob Sequence

randomized by

$$f(\underbrace{1, 0, 0, 1}_{\text{randomized by}}, \underbrace{\text{coin}, \text{coin}}_{\text{averaged over}})$$

averaged over

Doob Sequence

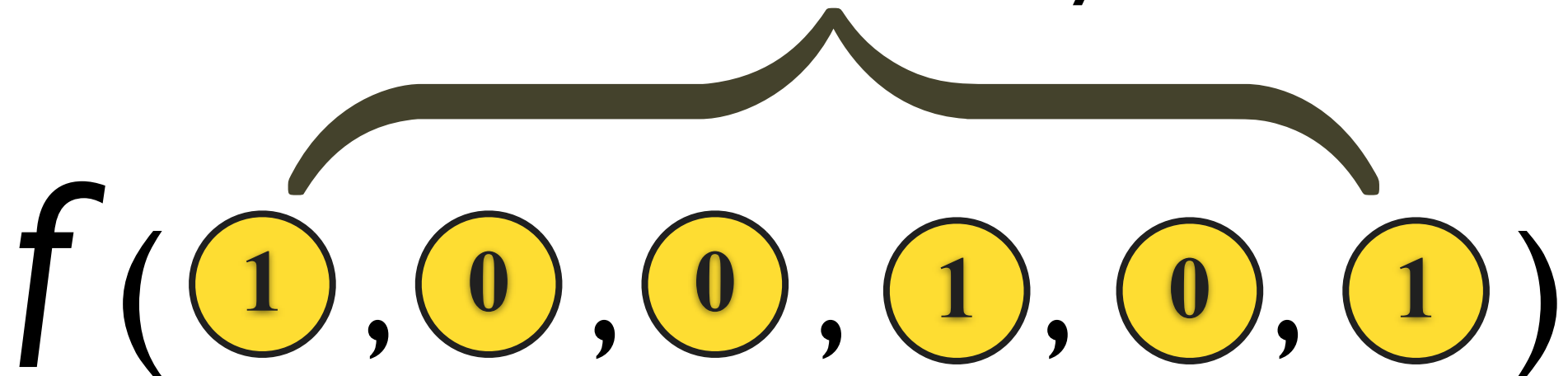
randomized by

$$f(\textcircled{1}, \textcircled{0}, \textcircled{0}, \textcircled{1}, \textcircled{0}, \textcircled{\$})$$

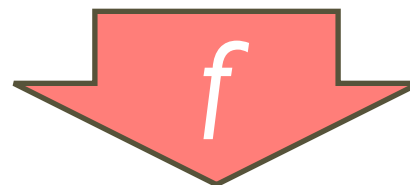
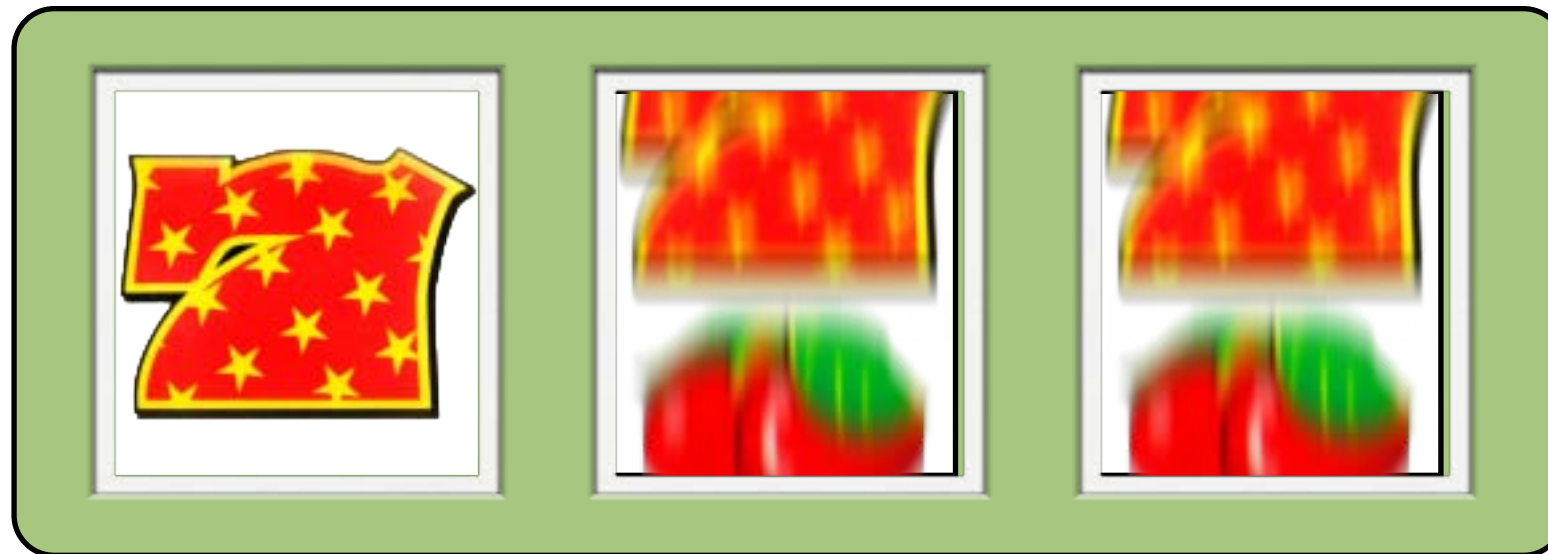
averaged over

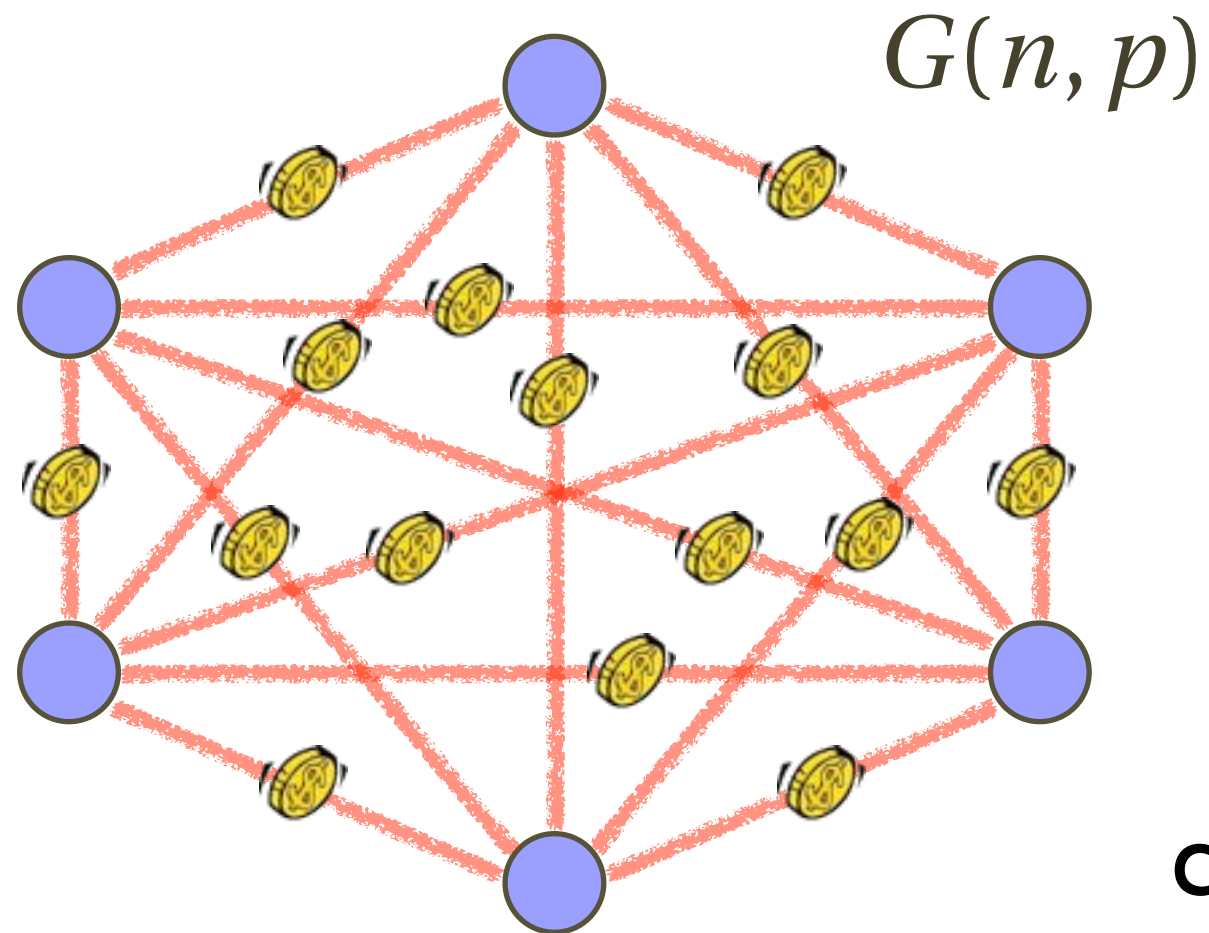
Doob Sequence

randomized by


$$f(1, 0, 0, 1, 0, 1)$$

Doob Martingale





Graph parameter:

$$f(G)$$

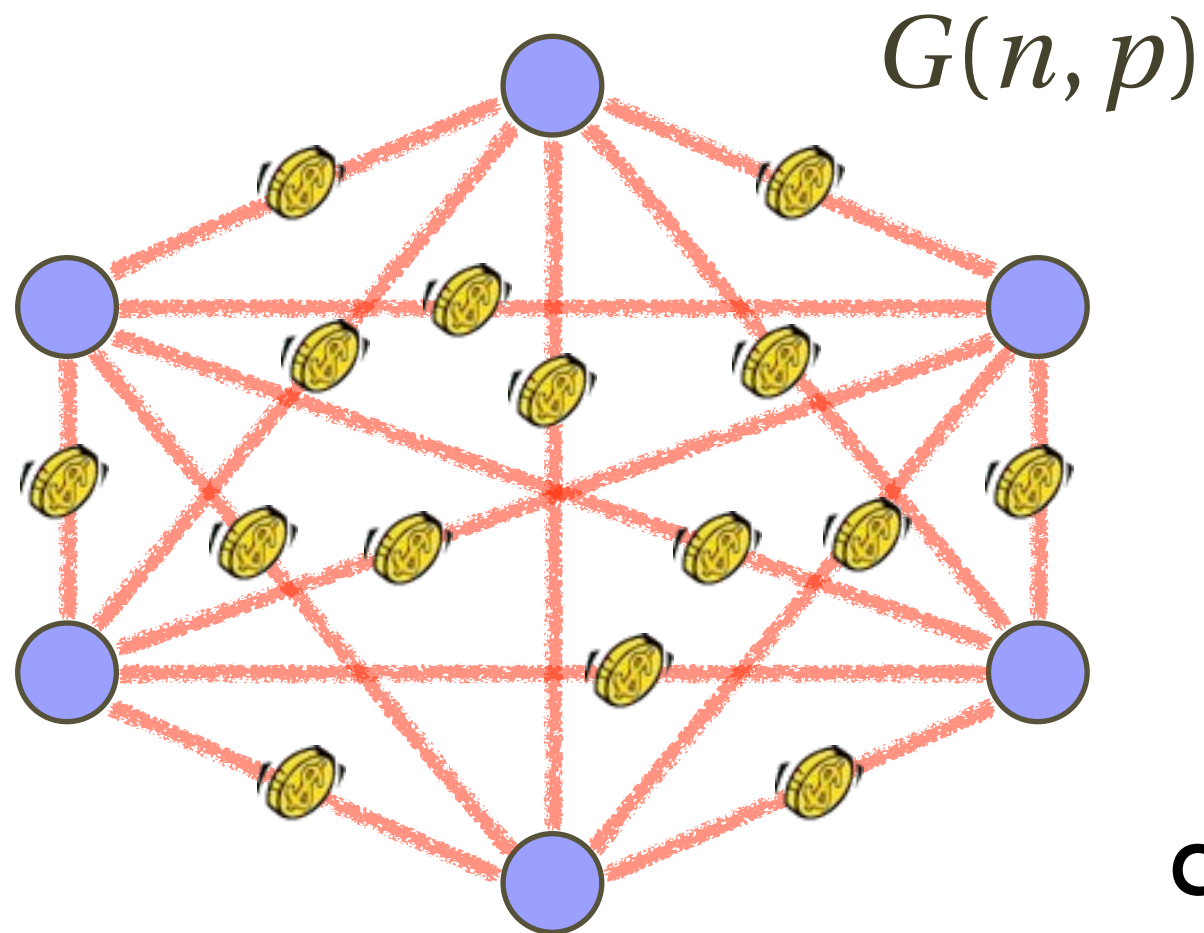
example: chromatic #,
components, diameter ...

numbering all vertex-pairs: $1, 2, 3, \dots, \binom{n}{2}$

$$I_j = \begin{cases} 1 & \text{edge } j \in G \\ 0 & \text{edge } j \notin G \end{cases}$$

$$Y_i = \mathbf{E}[f(G) \mid I_1, \dots, I_i]$$

$$Y_0 = \mathbf{E}[f(G)] \quad \text{-----} \rightarrow \quad Y_{\binom{n}{2}} = f(G)$$



Graph parameter:

$$f(G)$$

example: chromatic #,
components, diameter ...

numbering all vertices: $1, 2, 3, \dots, n$

X_i : **subgraph** of G **induced** by the first i vertices

$$Y_i = \mathbf{E}[f(G) \mid X_1, \dots, X_i]$$

$$Y_0 = \mathbf{E}[f(G)] \quad \text{-----} \rightarrow \quad Y_n = f(G)$$

Martingales induced by a random graph

- **Edge exposure** martingale:

I_j indicates the j th edge

$$Y_i = \mathbf{E}[f(G) \mid I_1, \dots, I_i]$$

- **Vertex exposure** martingale:

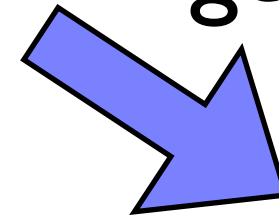
$$X_i = G([i])$$

$$Y_i = \mathbf{E}[f(G) \mid X_1, \dots, X_i]$$

martingale $X_0, X_1, X_2, \dots$

$$E[X_i | X_0, X_1, \dots, X_{i-1}] = X_{i-1}$$

generalization



martingale $Y_0, Y_1, Y_2, \dots$

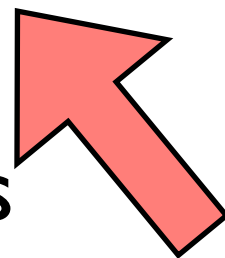
w.r.t. $X_0, X_1, X_2, \dots$

$$Y_i = f(X_0, X_1, \dots, X_i)$$

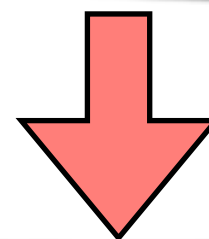
$$E[Y_i | X_0, X_1, \dots, X_{i-1}] = Y_{i-1}$$

edge-exposure martingale
vertex-exposure martingale

special cases
in random graphs



special case



Doob martingale

$$Y_i = E[f(X_0, X_1, \dots, X_n) | X_0, X_1, \dots, X_{i-1}]$$

Hoeffding's Inequality

Hoeffding's Inequality:

Let $X = \sum_{i=1}^n X_i$, where $X_1, \dots, X_n$ are independent random variables with

$$a_i \leq X_i \leq b_i.$$

Then

$$\Pr[|X - \mathbf{E}[X]| \geq t] \leq 2 \exp \left(- \frac{t^2}{2 \sum_{i=1}^n (b_i - a_i)^2} \right)$$

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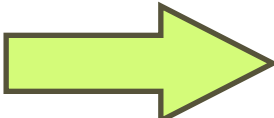
Then

$$\Pr[|X - \mathbf{E}[X]| \geq t] \leq 2 \exp\left(-\frac{t^2}{2 \sum_{i=1}^n (b_i - a_i)^2}\right)$$

Proof: X is a function (sum) of $X_1, \dots, X_n$.

Doob martingale: $Y_i = \mathbf{E}[X \mid X_1, \dots, X_i]$

$$|Y_i - Y_{i-1}| = |X_i - \mathbf{E}[X_i]| \leq b_i - a_i$$

Azuma  **Done!**

The Power of Doob + Azuma

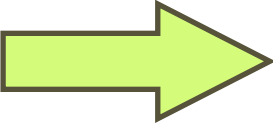
- For a function of (dependent) random variables: $f(X_1, \dots, X_n)$

- Doob martingale:

$$Y_i = \mathbf{E}[f(X_1, \dots, X_n) \mid X_1, \dots, X_i]$$

$$Y_0 = \mathbf{E}[f(X_1, \dots, X_n)] \quad Y_n = f(X_1, \dots, X_n)$$

- If the differences $|Y_i - Y_{i-1}|$ are bounded,

- Azuma  $|f(X_1, \dots, X_n) - \mathbf{E}[f(X_1, \dots, X_n)]|$ is small with high probability.
concentrated to its mean

The method of bounded differences:

Let $\mathbf{X} = (X_1, \dots, X_n)$ and let f be a function of $X_0, X_1, \dots, X_n$ satisfying that, for all $1 \leq i \leq n$,

$$|\mathbf{E}[f(\mathbf{X}) \mid X_1, \dots, X_i] - \mathbf{E}[f(\mathbf{X}) \mid X_1, \dots, X_{i-1}]| \leq c_i,$$

Then

$$\Pr \left[|f(\mathbf{X}) - \mathbf{E}[f(\mathbf{X})]| \geq t \right] \leq 2 \exp \left(- \frac{t^2}{2 \sum_{i=1}^n c_i^2} \right).$$

The method of bounded differences:

Let $X = (X_1, \dots, X_n)$ and let f be a function of $X_0, X_1, \dots, X_n$ satisfying that, for all $1 \leq i \leq n$,

$$|\mathbf{E}[f(X) \mid X_1, \dots, X_i] - \mathbf{E}[f(X) \mid X_1, \dots, X_{i-1}]| \leq c_i,$$

Y_i Y_{i-1}

Then

(Azuma)

$$\Pr \left[|\underbrace{f(X)}_{Y_n} - \underbrace{\mathbf{E}[f(X)]}_{Y_0}| \geq t \right] \leq 2 \exp \left(- \frac{t^2}{2 \sum_{i=1}^n c_i^2} \right).$$

Doob martingale: $Y_i = \mathbf{E}[f(X) \mid X_1, \dots, X_i]$

The method of bounded differences:

Let $\mathbf{X} = (X_1, \dots, X_n)$ and let f be a function of $X_0, X_1, \dots, X_n$ satisfying that, for all $1 \leq i \leq n$,

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Then

hard to check!

$$\Pr [|f(\mathbf{X}) - \mathbf{E}[f(\mathbf{X})]| \geq t] \leq 2 \exp \left(- \frac{t^2}{2 \sum_{i=1}^n c_i^2} \right).$$

~~$$|\mathbf{E}[f(\mathbf{X}) \mid X_1, \dots, X_i] - \mathbf{E}[f(\mathbf{X}) \mid X_1, \dots, X_{i-1}]| \leq c_i,$$~~

Lipschitz Condition:

$f(x_1, \dots, x_n)$ satisfies the **Lipschitz condition** with constants c_i , $1 \leq i \leq n$, if

$$\left| f(x_1, \dots, x_{i-1}, x_i, x_{i+1}, \dots, x_n) - f(x_1, \dots, x_{i-1}, y_i, x_{i+1}, \dots, x_n) \right| \leq c_i.$$

Average-case:

$$|\mathbf{E}[f(\mathbf{X}) \mid X_1, \dots, X_i] - \mathbf{E}[f(\mathbf{X}) \mid X_1, \dots, X_{i-1}]| \leq c_i,$$

Worst-case:

Lipschitz Condition:

$f(x_1, \dots, x_n)$ satisfies the **Lipschitz condition** with constants c_i , $1 \leq i \leq n$, if

$$\left| f(x_1, \dots, x_{i-1}, x_i, x_{i+1}, \dots, x_n) - f(x_1, \dots, x_{i-1}, y_i, x_{i+1}, \dots, x_n) \right| \leq c_i.$$

The method of bounded differences:

Let $\mathbf{X} = (X_1, \dots, X_n)$ be n independent random variables and let f be a function satisfying the Lipschitz condition with constants c_i , $1 \leq i \leq n$. Then

$$\Pr \left[|f(\mathbf{X}) - \mathbf{E}[f(\mathbf{X})]| \geq t \right] \leq 2 \exp \left(- \frac{t^2}{2 \sum_{i=1}^n c_i^2} \right).$$

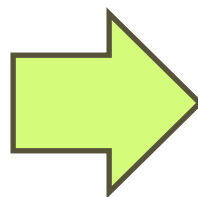
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Proof:

Lipschitz condition
+
independence



bounded averaged
differences

Applications of the Method

- A function f of **independent** random variable:

$$X_1, X_2, \dots, X_n$$

- **Lipschitz condition** of f :
 - changes of any variable makes little change to the value of f .
- $\Rightarrow f(X_1, X_2, \dots, X_n)$ is tightly concentrated to its mean.

Occupancy Problem

- m -balls-into- n -bins:
- number of empty bins?

$$X_i = \begin{cases} 1 & \text{bin } i \text{ is empty} \\ 0 & \text{otherwise} \end{cases}$$

empty bins: $X = \sum_{i=1}^n X_i$

$$\mathbf{E}[X] = \sum_{i=1}^n \mathbf{E}[X_i] = n \left(1 - \frac{1}{n}\right)^m$$

deviation: $\Pr[|X - \mathbf{E}[X]| \geq t] \leq ?$

X_i are
dependent

Occupancy Problem

- m -balls-into- n -bins:
- number of empty bins?

empty bins: X

deviation:

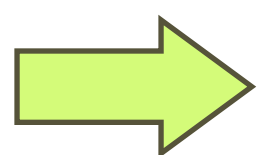
$$\Pr[|X - \mathbf{E}[X]| \geq t] \leq ?$$

Y_j : the bin of ball j (Independent!)

$$X = f(Y_1, \dots, Y_m) = |[n] - \{Y_1, \dots, Y_m\}|$$

Lipschitz:

changing any Y_j can change X for at most 1



$$\Pr[|X - \mathbf{E}[X]| \geq t\sqrt{m}] \leq 2e^{-t^2/2}$$

Pattern Matching

- a random string of length n ,
- a pattern of length k ,
- # of matched substrings?

alphabet Σ

$$|\Sigma| = m$$

a fixed
pattern: $\pi \in \Sigma^k$

uniform & **independent**: $X_1, \dots, X_n \in \Sigma$

Y : #substrings π in $(X_1, \dots, X_n)$

$$\mathbf{E}[Y] = (n - k + 1) \left(\frac{1}{m} \right)^k$$

Deviation?

Pattern Matching

- a random string of length n ,
- a pattern of length k ,
- # of matched substrings?

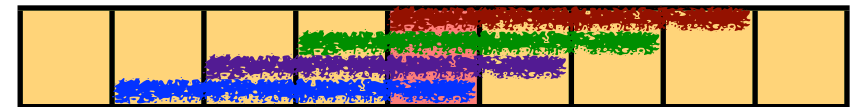
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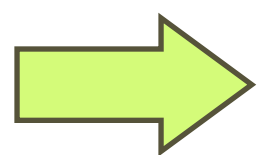
a fixed
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uniform & **independent:** $X_1, \dots, X_n \in \Sigma$

$$Y = f(X_1, \dots, X_n)$$

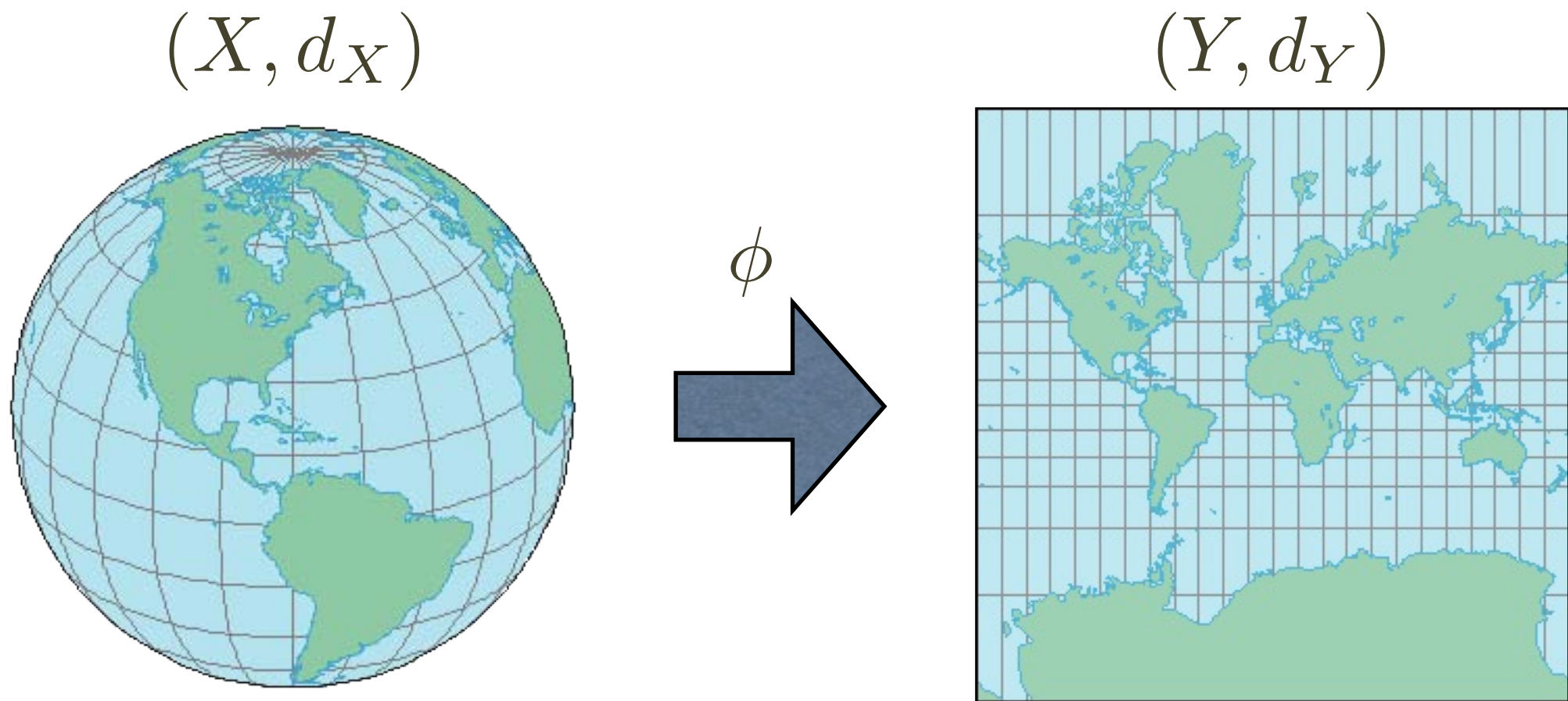


changing any X_i changes f for at most k



$$\Pr[|Y - \mathbf{E}[Y]| \geq tk\sqrt{n}] \leq 2e^{-t^2/2}$$

Metric Embedding



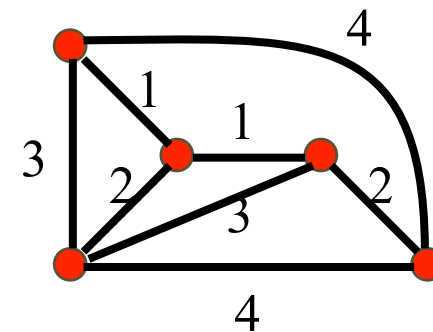
low-distortion: For a small $\alpha \geq 1$

$$\forall x_1, x_2 \in X, \quad \frac{1}{\alpha} d_X(x_1, x_2) \leq d_Y(\phi(x_1), \phi(x_2)) \leq \alpha d_X(x_1, x_2)$$

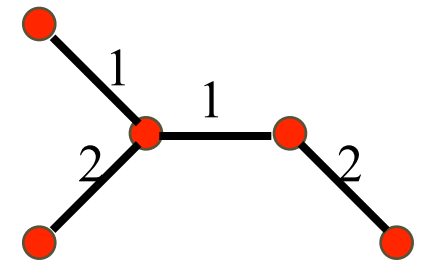
Why Embedding?

- Some problems are easier to solve in simple metrics:

- TSP in trees.



graph metric



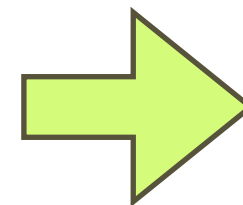
tree metric

- “Curse of dimensionality”

- proximity search;
- learning;
- due to volume explosion.



high dimension



low dimension

Dimension Reduction

In **Euclidian space**, it is always possible to embed a set of n points in **arbitrary** dimension to $O(\log n)$ dimension with constant distortion.

Johnson-Lindenstrauss Theorem:

For any $0 < \epsilon < 1$, for any set V of n points in $\mathbf{R}^d$, there is a map $\phi : \mathbf{R}^d \rightarrow \mathbf{R}^k$ with $k = O(\ln n)$, such that $\forall u, v \in V$,

$$(1 - \epsilon) \|u - v\|^2 \leq \|\phi(u) - \phi(v)\|^2 \leq (1 + \epsilon) \|u - v\|^2$$

Johnson-Lindenstrauss Theorem

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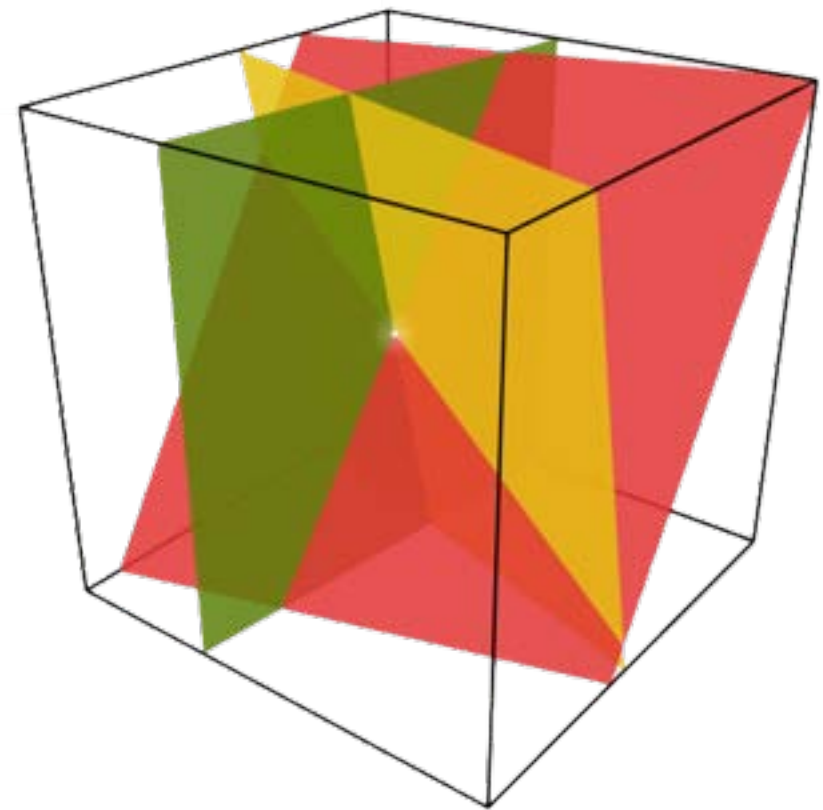
$$(1 - \epsilon) \|u - v\|^2 \leq \|\phi(u) - \phi(v)\|^2 \leq (1 + \epsilon) \|u - v\|^2$$

- $\phi(v) = Av$.
- A is a random projection matrix.

Random Projection

Random $k \times d$ matrix A :

- Projection onto a uniform random subspace.
(Johnson-Lindenstrauss)
(Dasgupta-Gupta)
- i.i.d. Gaussian entries.
(Indyk-Motiwani)
- i.i.d. $-1/+1$ entries.
(Achlioptas)



rows: $A_{1.}, A_{2.}, \dots, A_{k.}$

random orthogonal
unit vectors $\in \mathbb{R}^d$