

Randomized Algorithms

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Chernoff Bound

independent $X_1, X_2, \dots, X_n \in \{0, 1\}$

$$\text{let } X = \sum_{i=1}^n X_i$$

$t > 0$:

$$\Pr[X \geq \mathbf{E}[X] + t] \leq \exp\left(-\frac{2t^2}{n}\right)$$

$$\Pr[X \leq \mathbf{E}[X] - t] \leq \exp\left(-\frac{2t^2}{n}\right)$$

The method of bounded differences

Independent random variables: $X=(X_1, X_2, \dots, X_n)$.

$f(x_1, x_2, \dots, x_n)$ satisfies the Lipschitz condition:

$$|f(x_1, \dots, x_n) - f(x_1, \dots, x_{i-1}, y_i, x_{i+1}, \dots, x_n)| \leq c_i$$

for arbitrary possible values x_1, \dots, x_n, y_i .

$t > 0$:

$$\Pr[f(\mathbf{X}) \geq \mathbf{E}[f(\mathbf{X})] + t] \leq \exp\left(-\frac{2t^2}{\sum_{i=1}^n c_i^2}\right)$$

$$\Pr[f(\mathbf{X}) \leq \mathbf{E}[f(\mathbf{X})] - t] \leq \exp\left(-\frac{2t^2}{\sum_{i=1}^n c_i^2}\right)$$

Martingale

Definition:

A sequence of random variables X_0, X_1, \dots is a **martingale** if for all $i > 0$,

$$\mathbf{E}[X_i | X_0, \dots, X_{i-1}] = X_{i-1}$$

What does this mean?

Azuma's Inequality:

Let X_0, X_1, \dots be a martingale such that, for all $k \geq 1$,

$$|X_k - X_{k-1}| \leq c_k,$$

Then

$$\Pr [|X_n - X_0| \geq t] \leq 2 \exp \left(-\frac{t^2}{2 \sum_{k=1}^n c_k^2} \right).$$

Generalization

Definition:

Y_0, Y_1, \dots is a martingale with respect to X_0, X_1, \dots if, for all $i \geq 0$,

- Y_i is a function of X_0, X_1, \dots, X_i ;
- $\mathbb{E}[Y_{i+1} | X_0, \dots, X_i] = Y_i$.

Definition:

Y_0, Y_1, \dots is a martingale with respect to X_0, X_1, \dots if, for all $i \geq 0$,

- Y_i is a function of X_0, X_1, \dots, X_i ;
- $\mathbb{E}[Y_{i+1} | X_0, \dots, X_i] = Y_i$.

- Betting on a fair game;
- X_i : win/loss of the i -th bet;
- Y_i : wealth after the i -th bet -- Martingale (fair game)

Azuma's Inequality (general version):

Let Y_0, Y_1, \dots be a martingale with respect to X_0, X_1, \dots such that, for all $k \geq 1$,

$$|Y_k - Y_{k-1}| \leq c_k,$$

Then

$$\Pr [|Y_n - Y_0| \geq t] \leq 2 \exp \left(-\frac{t^2}{2 \sum_{k=1}^n c_k^2} \right).$$

Doob Sequence

Definition (Doob sequence):

The Doob sequence of a function f with respect to a sequence X_1, \dots, X_n is

$$Y_i = \mathbf{E}[f(X_1, \dots, X_n) \mid X_1, \dots, X_i]$$

$$Y_0 = \mathbf{E}[f(X_1, \dots, X_n)] \xrightarrow{\text{-----}} Y_n = f(X_1, \dots, X_n)$$

Doob Sequence

$$f(\langle \text{coin}, \text{coin}, \text{coin}, \text{coin}, \text{coin}, \text{coin} \rangle)$$

averaged over

$$\mathbf{E}[f] = Y_0,$$

Doob Sequence

randomized by

$$f(1, \text{coin}, \text{coin}, \text{coin}, \text{coin}, \text{coin}, \text{coin})$$



averaged over

$$\mathbb{E}[f] = Y_0, Y_1,$$

Doob Sequence

randomized by

$$f(1, 0, \text{heads}, \text{tails}, \text{heads}, \text{tails})$$

averaged over

$$\mathbb{E}[f] = Y_0, \quad Y_1, \quad Y_2,$$

Doob Sequence

randomized by

$$f(1, 0, 0, Y_0, Y_1, Y_2, Y_3)$$

averaged over

$$\mathbb{E}[f] = Y_0, Y_1, Y_2, Y_3,$$

Doob Sequence

randomized by

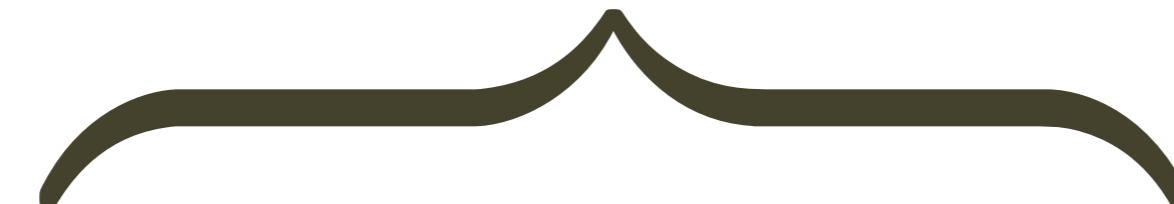
$$f(1, 0, 0, 1, Y_0, Y_1)$$

averaged over

$$\mathbb{E}[f] = Y_0, Y_1, Y_2, Y_3, Y_4,$$

Doob Sequence

randomized by


$$f((1, 0, 0, 1, 0, ?))$$

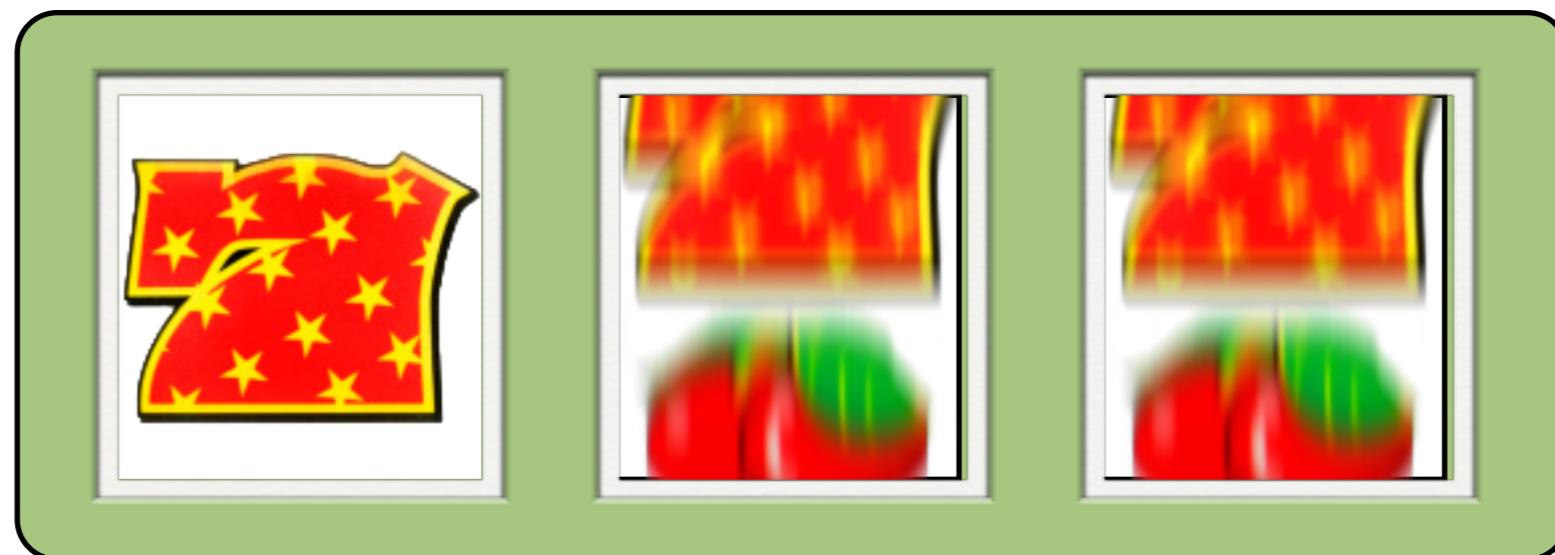
averaged over

$$\mathbb{E}[f] = Y_0, \quad Y_1, \quad Y_2, \quad Y_3, \quad Y_4, \quad Y_5,$$

Doob Sequence

$$\mathbf{E}[f] = Y_0, \quad Y_1, \quad Y_2, \quad Y_3, \quad Y_4, \quad Y_5, \quad Y_6 \quad = f$$

Doob Martingale



Doob sequence:

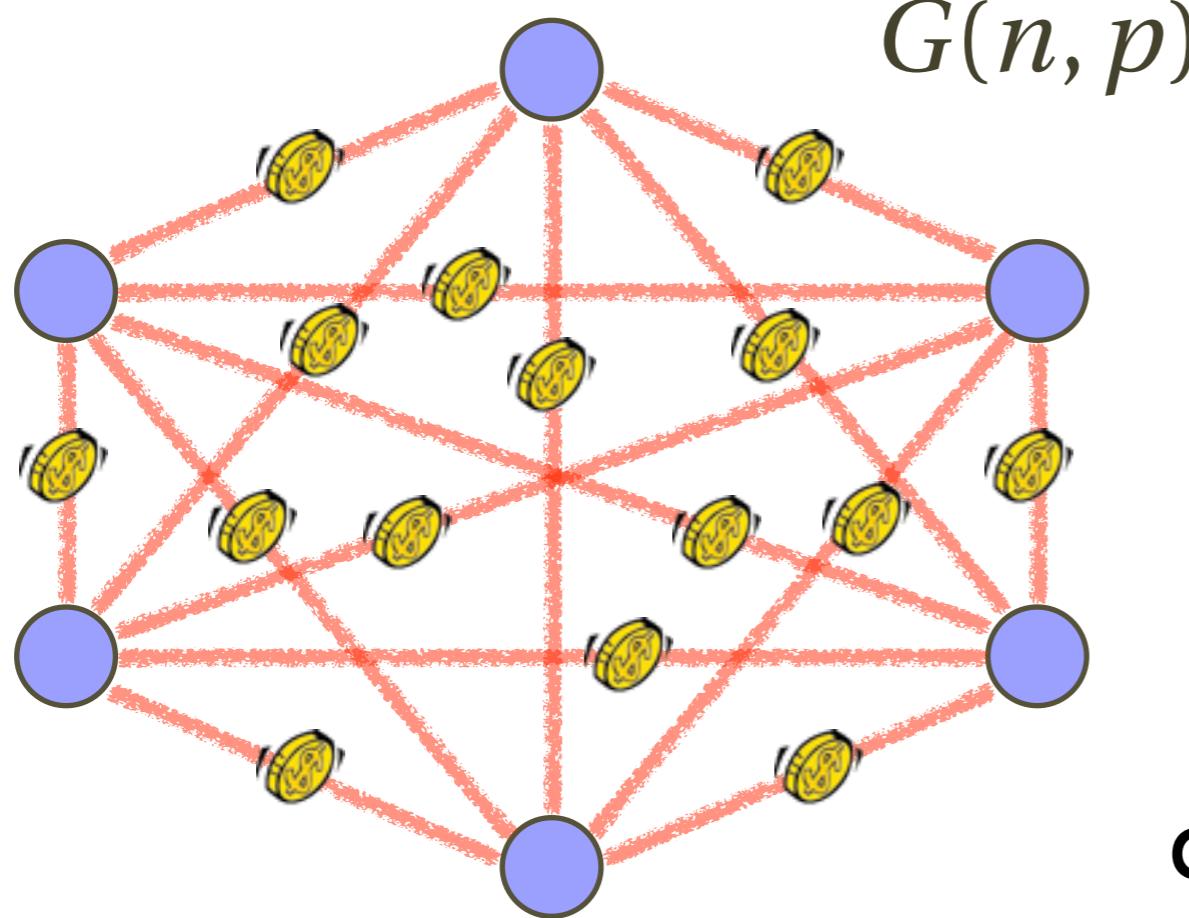
$$Y_i = \mathbf{E}[f(X_1, \dots, X_n) \mid X_1, \dots, X_i]$$

Doob sequence is a martingale:

$$\mathbf{E}[Y_i \mid X_1, \dots, X_{i-1}] = Y_{i-1}$$

Proof:

$$\begin{aligned} & \mathbf{E}[Y_i \mid X_1, \dots, X_{i-1}] \\ &= \mathbf{E}[\mathbf{E}[f(X_1, \dots, X_n) \mid X_1, \dots, X_i] \mid X_1, \dots, X_{i-1}] \\ &= \mathbf{E}[f(X_1, \dots, X_n) \mid X_1, \dots, X_{i-1}] \\ &= Y_{i-1} \end{aligned}$$



Graph parameter:

$$f(G)$$

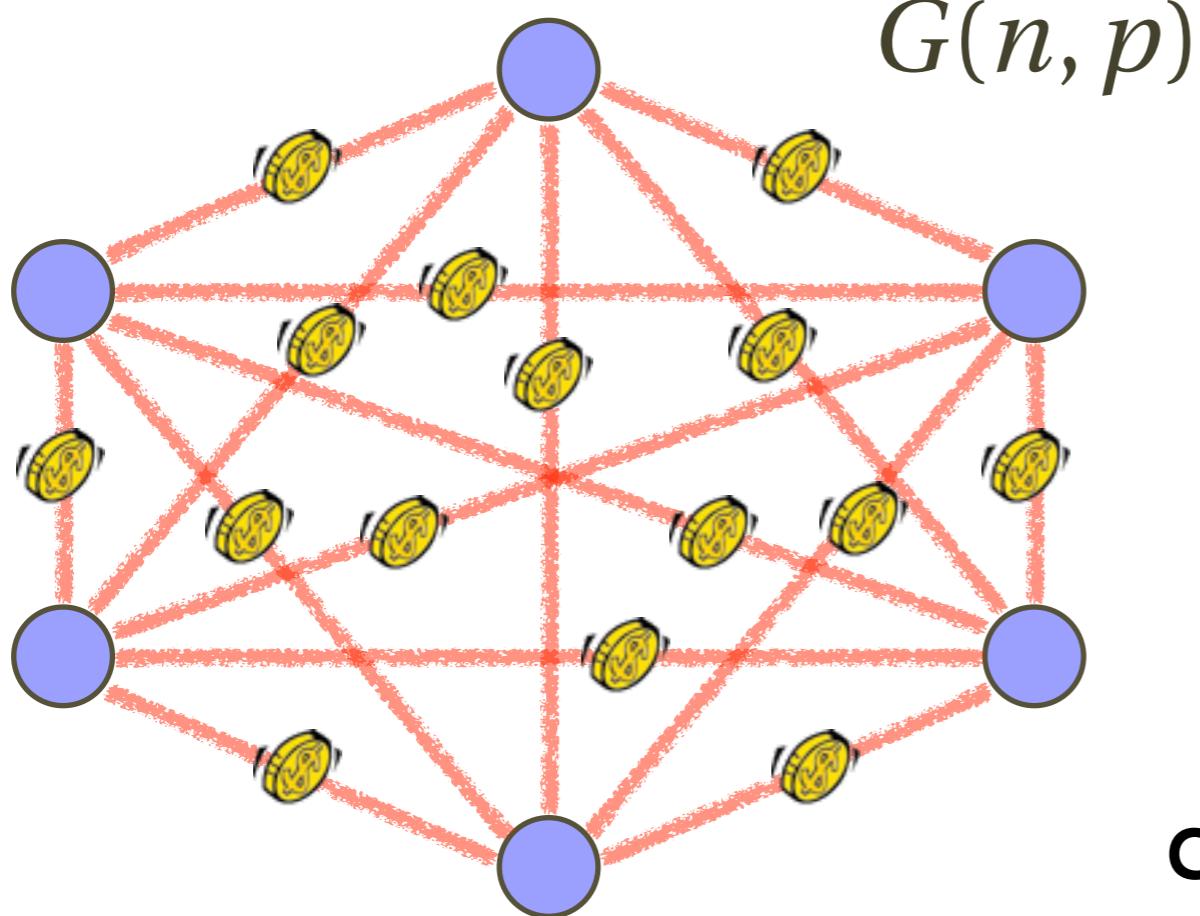
example: chromatic #,
components, diameter ...

numbering all vertex-pairs: $1, 2, 3, \dots, \binom{n}{2}$

$$I_j = \begin{cases} 1 & \text{edge } j \in G \\ 0 & \text{edge } j \notin G \end{cases}$$

$$Y_i = \mathbf{E}[f(G) \mid I_1, \dots, I_i]$$

$$Y_0 = \mathbf{E}[f(G)] \quad \xrightarrow{\hspace{1cm}} \quad Y_{\binom{n}{2}} = f(G)$$



Graph parameter:

$$f(G)$$

example: chromatic #,
components, diameter ...

numbering all vertices: $1, 2, 3, \dots, n$

X_i : **subgraph** of G induced by the first i vertices

$$Y_i = \mathbf{E}[f(G) \mid X_1, \dots, X_i]$$

$$Y_0 = \mathbf{E}[f(G)] \quad \xrightarrow{\hspace{1cm}} \quad Y_n = f(G)$$

Martingales induced by a random graph

- **Edge exposure martingale:**

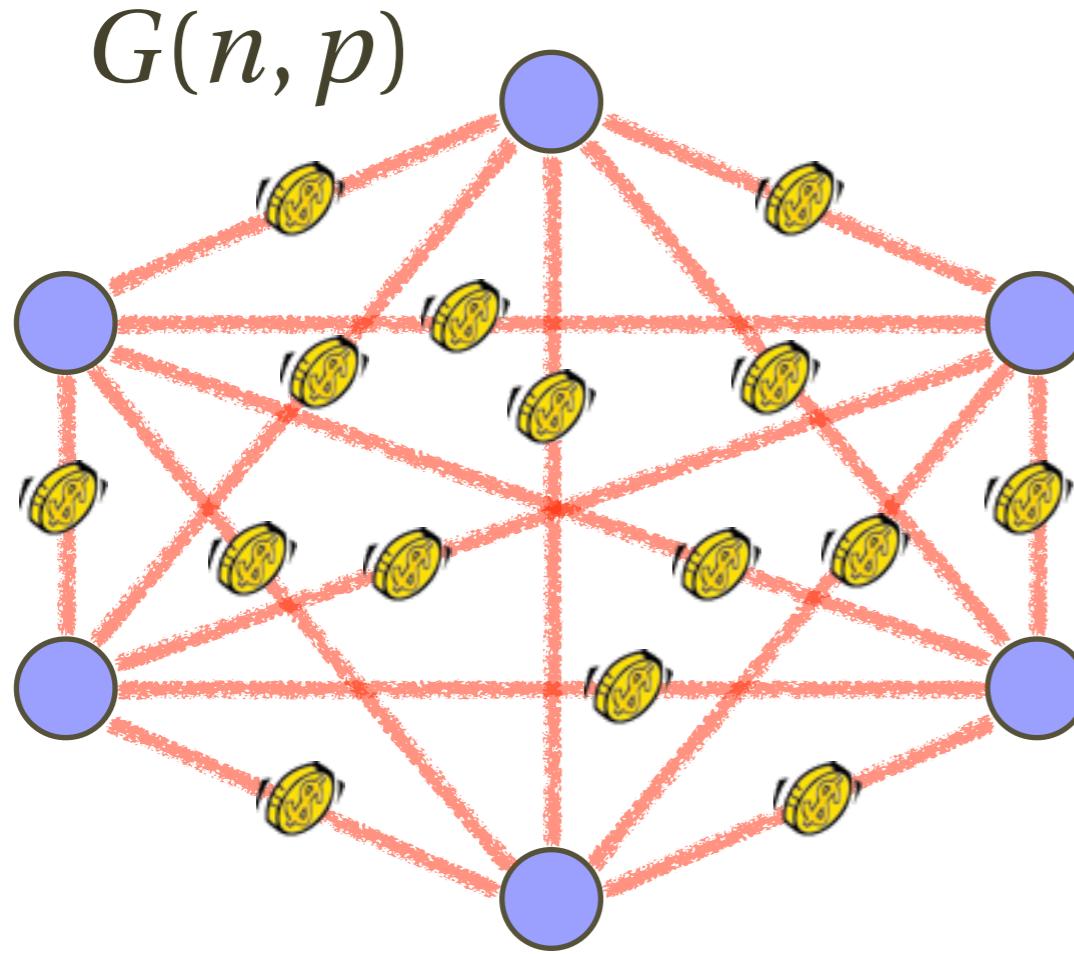
I_j indicates the j th edge

$$Y_i = \mathbf{E}[f(G) \mid I_1, \dots, I_i]$$

- **Vertex exposure martingale:**

$$X_i = G([i])$$

$$Y_i = \mathbf{E}[f(G) \mid X_1, \dots, X_i]$$



chromatic number:
 $\chi(G)$

the **smallest** number of colors to **properly color** G

X_i : **subgraph** of G **induced** by the first i vertices

$$Y_i = \mathbb{E}[\chi(G) \mid X_1, \dots, X_i]$$

Doob martingale: Y_0, Y_1, \dots, Y_n

$$Y_0 = \mathbb{E}[\chi(G)]$$

$$Y_n = \chi(G)$$

chromatic number: $\chi(G)$

X_i : **subgraph of G induced by the first i vertices**

$$Y_i = \mathbf{E}[\chi(G) \mid X_1, \dots, X_i]$$

Doob martingale: Y_0, Y_1, \dots, Y_n

$$Y_0 = \mathbf{E}[\chi(G)] \qquad \qquad Y_n = \chi(G)$$

Observation:

a vertex can always be given a new color.

$$|Y_i - Y_{i-1}| \leq 1$$

Azuma's Inequality (general version):

Let Y_0, Y_1, \dots be a martingale with respect to X_0, X_1, \dots such that, for all $k \geq 1$,

$$|Y_k - Y_{k-1}| \leq c_k,$$

Then

$$\Pr[|Y_n - Y_0| \geq t] \leq 2 \exp\left(-\frac{t^2}{2 \sum_{k=1}^n c_k^2}\right).$$

Observation:

a vertex can always be given a new color.

$$|Y_i - Y_{i-1}| \leq 1$$

$$\Pr[|\chi(G) - \mathbf{E}[\chi(G)]| \geq t\sqrt{n}]$$

$$= \Pr[|Y_n - Y_0| \geq t\sqrt{n}] \leq 2e^{-t^2/2}$$

Tight Concentration of Chromatic Number

Theorem [Shamir & Spencer (1987)]:

Let $G \sim G(n, p)$. Then

$$\Pr[|\chi(G) - \mathbb{E}[\chi(G)]| \geq t\sqrt{n}] \leq 2e^{-t^2/2}.$$

martingale X_0, X_1, X_2, \dots

$$E[X_i | X_0, X_1, \dots, X_{i-1}] = X_{i-1}$$

generalization

edge-exposure martingale
vertex-exposure martingale

martingale Y_0, Y_1, Y_2, \dots

w.r.t. X_0, X_1, X_2, \dots

$$Y_i = f(X_0, X_1, \dots, X_i)$$

$$E[Y_i | X_0, X_1, \dots, X_{i-1}] = Y_{i-1}$$

special cases
in random graphs

Doob martingale

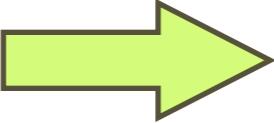
$$Y_i = E[f(X_0, X_1, \dots, X_i) | X_0, X_1, \dots, X_{i-1}]$$

The Power of Doob + Azuma

- For a function of (dependent) random variables: $f(X_1, \dots, X_n)$
- Doob martingale:

$$Y_i = \mathbf{E}[f(X_1, \dots, X_n) \mid X_1, \dots, X_i]$$

$$Y_0 = \mathbf{E}[f(X_1, \dots, X_n)] \quad Y_n = f(X_1, \dots, X_n)$$

- If the differences $|Y_i - Y_{i-1}|$ are bounded,
- Azuma  $|Y_n - Y_0|$ is small whp.

$f(X_1, \dots, X_n)$ is tightly concentrated to its mean

The method of bounded differences:

Let $X = (X_1, \dots, X_n)$ and let f be a function of X_0, X_1, \dots, X_n satisfying that, for all $1 \leq i \leq n$,

$$|\mathbf{E}[f(X) | X_1, \dots, X_i] - \mathbf{E}[f(X) | X_1, \dots, X_{i-1}]| \leq c_i,$$

Then

$$\Pr [|f(X) - \mathbf{E}[f(X)]| \geq t] \leq 2 \exp\left(-\frac{t^2}{2 \sum_{i=1}^n c_i^2}\right).$$

The method of bounded differences:

Let $X = (X_1, \dots, X_n)$ and let f be a function of X_0, X_1, \dots, X_n satisfying that, for all $1 \leq i \leq n$,

$$|\mathbb{E}[f(X) \mid X_1, \dots, X_i] - \mathbb{E}[f(X) \mid X_1, \dots, X_{i-1}]| \leq c_i,$$

$$\Pr \left[|f(X) - \mathbb{E}[f(X)]| \geq t \right] \leq 2 \exp \left(-\frac{t^2}{2 \sum_{i=1}^n c_i^2} \right).$$

$$\begin{matrix} Y_i & & Y_{i-1} \\ & \vdots & \\ Y_n & & Y_0 \end{matrix}$$

Then

(Azuma) $\Pr \left[|f(X) - \mathbb{E}[f(X)]| \geq t \right] \leq 2 \exp \left(-\frac{t^2}{2 \sum_{i=1}^n c_i^2} \right).$

Doob martingale: $Y_i = \mathbb{E}[f(X) \mid X_1, \dots, X_i]$

The method of bounded differences:

Let $X = (X_1, \dots, X_n)$ and let f be a function of X_0, X_1, \dots, X_n satisfying that, for all $1 \leq i \leq n$,

$$|\mathbf{E}[f(X) | X_1, \dots, X_i] - \mathbf{E}[f(X) | X_1, \dots, X_{i-1}]| \leq c_i,$$

Then

hard to check!

$$\Pr [|f(X) - \mathbf{E}[f(X)]| \geq t] \leq 2 \exp \left(-\frac{t^2}{2 \sum_{i=1}^n c_i^2} \right).$$

$$|\mathbf{E}[f(X) | X_1, \dots, X_i] - \mathbf{E}[f(X) | X_1, \dots, X_{i-1}]| \leq c_i,$$

Lipschitz Condition:

$f(x_1, \dots, x_n)$ satisfies the **Lipschitz condition** with constants c_i , $1 \leq i \leq n$, if

$$\left| f(x_1, \dots, x_{i-1}, x_i, x_{i+1}, \dots, x_n) \right. \\ \left. - f(x_1, \dots, x_{i-1}, y_i, x_{i+1}, \dots, x_n) \right| \leq c_i.$$

Average-case:

$$|\mathbb{E}[f(X) \mid X_1, \dots, X_i] - \mathbb{E}[f(X) \mid X_1, \dots, X_{i-1}]| \leq c_i,$$

Worst-case:

Lipschitz Condition:

$f(x_1, \dots, x_n)$ satisfies the **Lipschitz condition** with constants c_i , $1 \leq i \leq n$, if

$$\left| f(x_1, \dots, x_{i-1}, x_i, x_{i+1}, \dots, x_n) \right. \\ \left. - f(x_1, \dots, x_{i-1}, y_i, x_{i+1}, \dots, x_n) \right| \leq c_i.$$

The method of bounded differences:

Let $X = (X_1, \dots, X_n)$ be n independent random variables and let f be a function satisfying the Lipschitz condition with constants c_i , $1 \leq i \leq n$. Then

$$\Pr [|f(X) - \mathbf{E}[f(X)]| \geq t] \leq 2 \exp \left(-\frac{t^2}{2 \sum_{i=1}^n c_i^2} \right).$$

The method of bounded differences:

Let $X = (X_1, \dots, X_n)$ be n independent random variables and let f be a function satisfying the Lipschitz condition with constants c_i , $1 \leq i \leq n$. Then

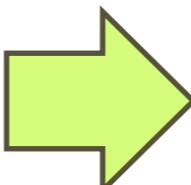
$$\Pr [|f(X) - \mathbb{E}[f(X)]| \geq t] \leq 2 \exp \left(-\frac{t^2}{2 \sum_{i=1}^n c_i^2} \right).$$

Proof:

Lipschitz condition

+

independence



bounded averaged
differences

Occupancy Problem

- m -balls-into- n -bins:
- number of empty bins?

$$X_i = \begin{cases} 1 & \text{bin } i \text{ is empty} \\ 0 & \text{otherwise} \end{cases}$$

empty bins: $X = \sum_{i=1}^n X_i$

$$\mathbb{E}[X] = \sum_{i=1}^n \mathbb{E}[X_i] = n \left(1 - \frac{1}{n}\right)^m$$

deviation: $\Pr[|X - \mathbb{E}[X]| \geq t] \leq ?$

X_i are
dependent

Occupancy Problem

- m -balls-into- n -bins:
- number of empty bins?

empty bins: X

deviation:

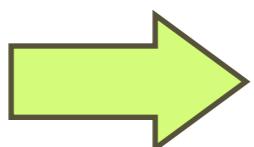
$$\Pr[|X - \mathbb{E}[X]| \geq t] \leq ?$$

Y_j : the bin of ball j (**Independent!**)

$$X = f(Y_1, \dots, Y_m) = |[n] - \{Y_1, \dots, Y_m\}|$$

Lipschitz:

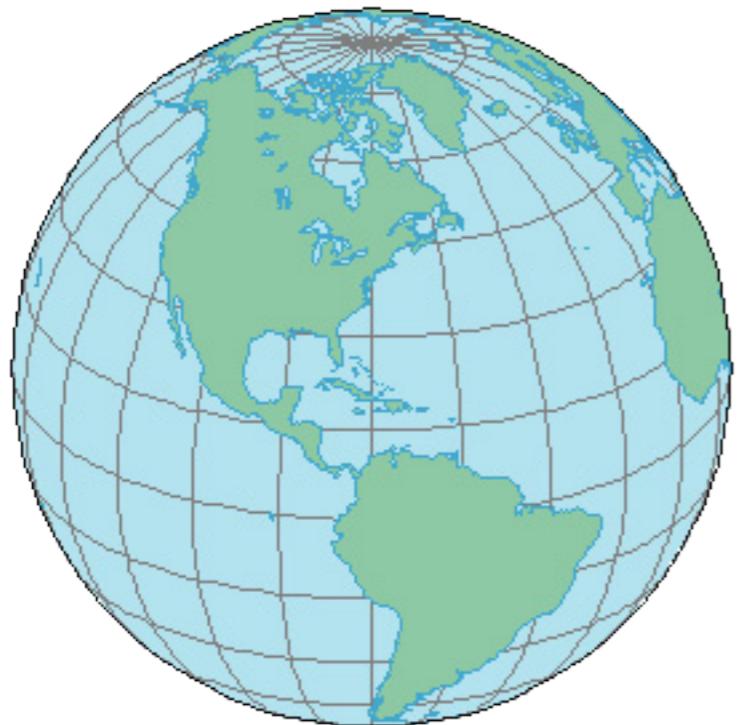
changing any Y_j can change X for at most 1



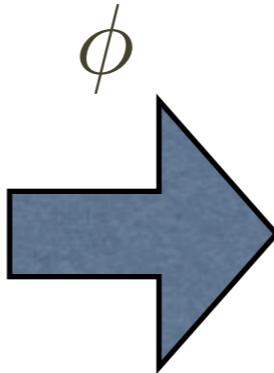
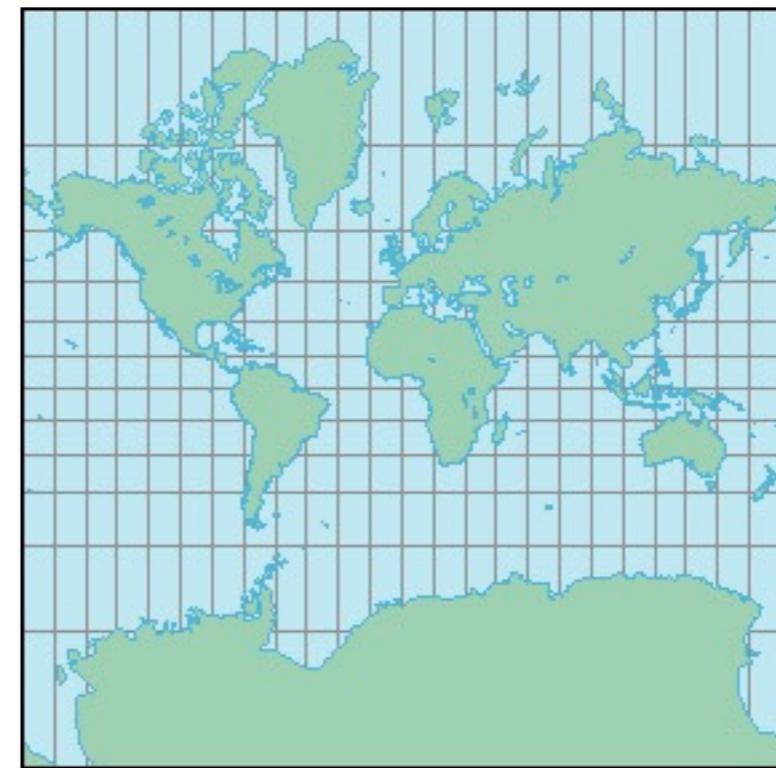
$$\Pr[|X - \mathbb{E}[X]| \geq t\sqrt{m}] \leq 2e^{-t^2/2}$$

Metric Embedding

(X, d_X)



(Y, d_Y)



low-distortion: For a small $\alpha \geq 1$

$$\forall x_1, x_2 \in X, \quad \frac{1}{\alpha} d_X(x_1, x_2) \leq d_Y(\phi(x_1), \phi(x_2)) \leq \alpha d_X(x_1, x_2)$$

Dimension Reduction

In **Euclidian space**, it is always possible to embed a set of n points in **arbitrary** dimension to $O(\log n)$ dimension with constant distortion.

Johnson-Lindenstrauss Theorem:

For any $0 < \epsilon < 1$, for any set V of n points in \mathbf{R}^d , there is a map $\phi : \mathbf{R}^d \rightarrow \mathbf{R}^k$ with $k = O(\ln n)$, such that $\forall u, v \in V$,

$$(1 - \epsilon) \|u - v\|^2 \leq \|\phi(u) - \phi(v)\|^2 \leq (1 + \epsilon) \|u - v\|^2$$

Johnson-Lindenstrauss Theorem

Johnson-Lindenstrauss Theorem:

For any $0 < \epsilon < 1$, for any set V of n points in \mathbf{R}^d ,
there is a map $\phi : \mathbf{R}^d \rightarrow \mathbf{R}^k$ with $k = O(\ln n)$,
such that $\forall u, v \in V$,

$$(1 - \epsilon) \|u - v\|^2 \leq \|\phi(u) - \phi(v)\|^2 \leq (1 + \epsilon) \|u - v\|^2$$

- $\phi(v) = Av$.
- A is a random projection matrix.

Random Projection

Random $k \times d$ matrix A :

- Projection onto a uniform random subspace.

(Johnson-Lindenstrauss)

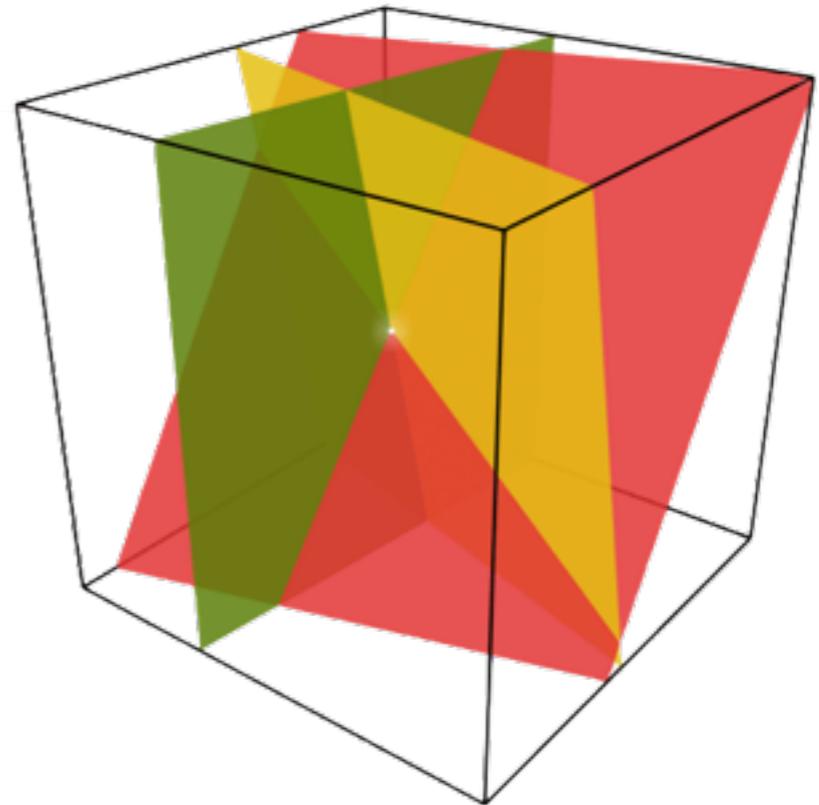
(Dasgupta-Gupta)

- i.i.d. Gaussian entries.

(Indyk-Motiwani)

- i.i.d. -1/+1 entries.

(Achlioptas)



rows: $A_{1\cdot}, A_{2\cdot}, \dots, A_{k\cdot}$.

random orthogonal
unit vectors $\in \mathbb{R}^d$

Johnson-Lindenstrauss Theorem

- Let V be a set of n points in \mathbf{R}^d .
- For some $k = O(\ln n)$.
- Let A be a random $k \times d$ matrix that projects onto a uniform random k -dimensional subspace.

W.h.p., $\forall u, v \in V$,

$$\frac{\|u - v\|^2}{2} \leq \left\| \sqrt{\frac{d}{k}} A u - \sqrt{\frac{d}{k}} A v \right\|^2 \leq \frac{3\|u - v\|^2}{2}$$

(for the case $\varepsilon=1/2$)

the projection

A : projection onto a uniform
random k -subspace

W.h.p., $\forall u, v \in V$,

$$\left\| \sqrt{\frac{d}{k}}Au - \sqrt{\frac{d}{k}}Av \right\|^2 \leq \frac{3\|u - v\|^2}{2}$$

$$\left\| \sqrt{\frac{d}{k}}Au - \sqrt{\frac{d}{k}}Av \right\|^2 \geq \frac{\|u - v\|^2}{2}$$

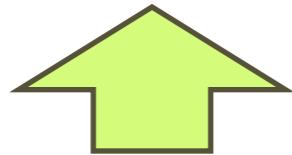
Step I:

Reduce to unit vectors

A : projection onto a uniform random k -subspace

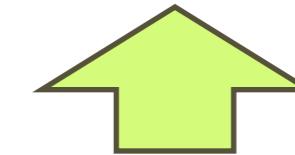
$O(n^2)$ pairs
W.h.p., $\forall u, v \in V$,

$$\left\| \sqrt{\frac{d}{k}} Au - \sqrt{\frac{d}{k}} Av \right\|^2 \leq \frac{3\|u - v\|^2}{2}$$



$$\left\| A \frac{(u - v)}{\|u - v\|} \right\|^2 \leq \frac{3k}{2d}$$

$$\left\| \sqrt{\frac{d}{k}} Au - \sqrt{\frac{d}{k}} Av \right\|^2 \geq \frac{\|u - v\|^2}{2}$$



$$\left\| A \frac{(u - v)}{\|u - v\|} \right\|^2 \geq \frac{k}{2d}$$

treat $(u - v)$ as a vector, $\frac{u-v}{\|u-v\|}$ is a unit vector

New goal:

\forall unit vector u , with probability $o(\frac{1}{n^2})$

$$\|Au\|^2 > \frac{3k}{2d} \quad \text{or} \quad \|Au\|^2 < \frac{k}{2d}$$

A : projection onto a uniform random k -subspace

\forall unit vector u , with probability $o(\frac{1}{n^2})$

$$\|Au\|^2 > \frac{3k}{2d} \quad \text{or} \quad \|Au\|^2 < \frac{k}{2d}$$

Step II:

Random projection of fixed unit vector

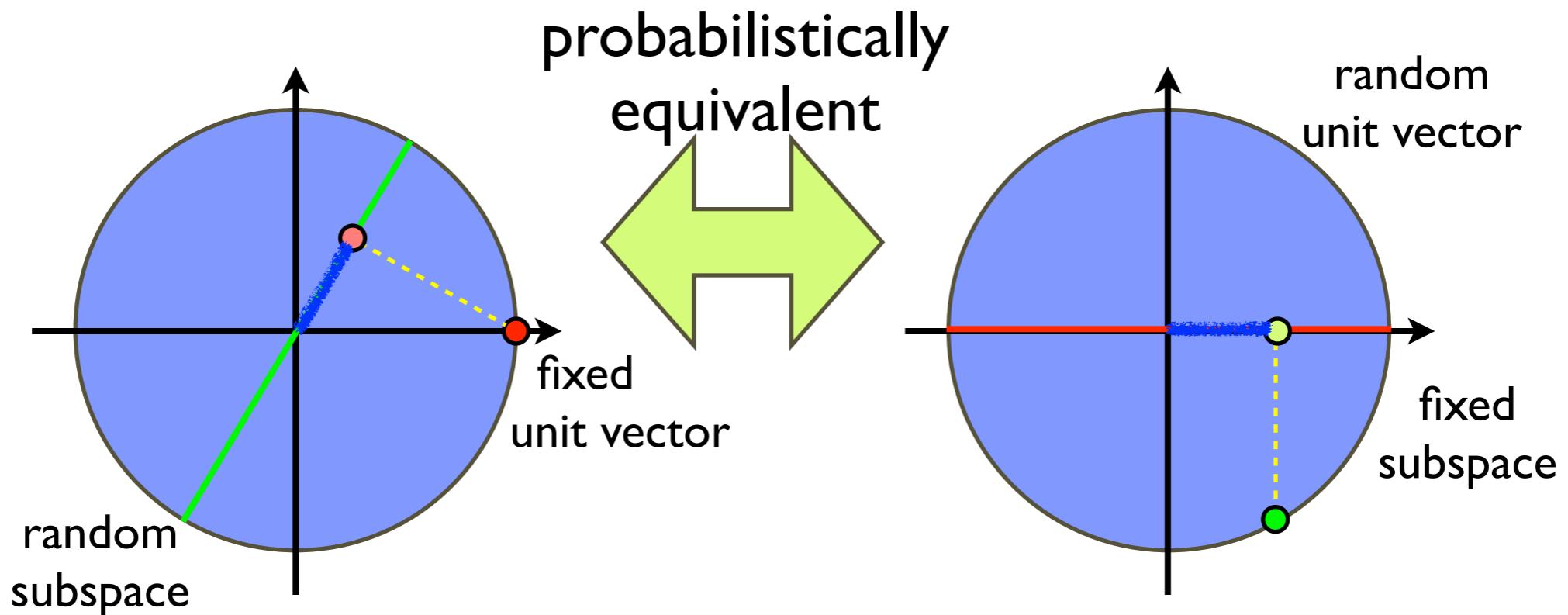


Fixed projection of random unit vector

A : projection onto a uniform random k -subspace

\forall unit vector u , with probability $o\left(\frac{1}{n^2}\right)$

$$\|Au\|^2 > \frac{3k}{2d} \quad \text{or} \quad \|Au\|^2 < \frac{k}{2d}$$



“inner-products are invariant under rotations”

A : projection onto a uniform random k -subspace

\forall unit vector u , with probability $o(\frac{1}{n^2})$

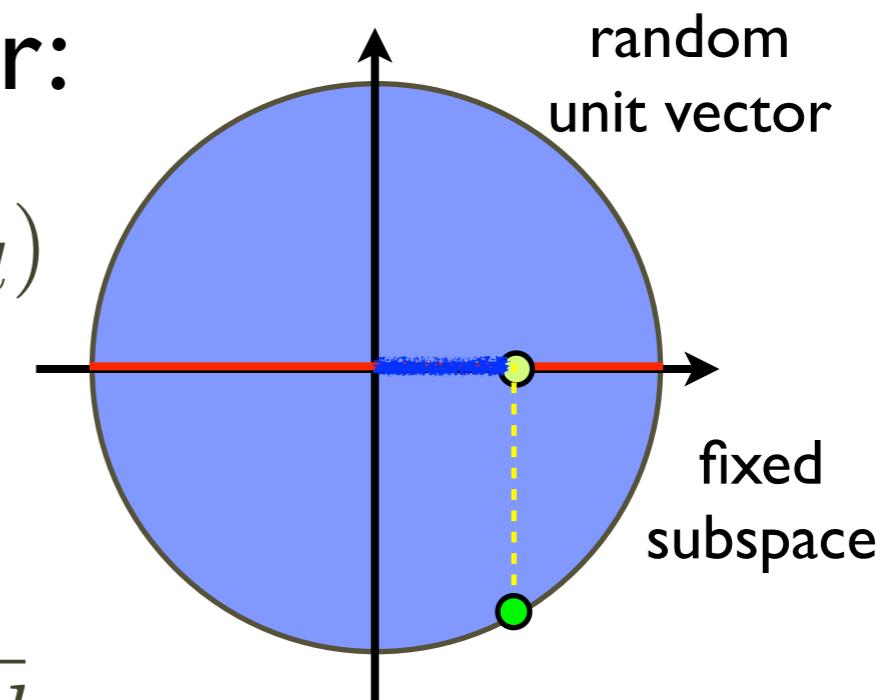
$$\|Au\|^2 > \frac{3k}{2d} \quad \text{or} \quad \|Au\|^2 < \frac{k}{2d}$$

uniform random **unit** vector:

$$Y = (Y_1, \dots, Y_k, Y_{k+1}, \dots, Y_d)$$

$$Z = (Y_1, \dots, Y_k)$$

$$\|Z\|^2 > \frac{3k}{2d} \quad \text{or} \quad \|Z\|^2 < \frac{k}{2d}$$



uniform random unit vector:

$$Y = (Y_1, \dots, Y_k, Y_{k+1}, \dots, Y_d)$$

Let $Z = (Y_1, \dots, Y_k)$ with probability $o(\frac{1}{n^2})$

$$\|Z\|^2 > \frac{3k}{2d} \quad \text{or} \quad \|Z\|^2 < \frac{k}{2d}$$

Looks clean, however . . .

- $\|Z\|^2 = Y_1^2 + Y_2^2 + \cdots + Y_k^2$
- Y_i 's are dependent for uniform unit vector Y .

uniform random **unit vector:**

$$Y = (Y_1, \dots, Y_k, Y_{k+1}, \dots, Y_d)$$

Let $Z = (Y_1, \dots, Y_k)$ with probability $o(\frac{1}{n^2})$

$$\|Z\|^2 > \frac{3k}{2d} \quad \text{or} \quad \|Z\|^2 < \frac{k}{2d}$$

Step III:
Bound $\|Z\|^2$.

uniform random **unit vector:**

$$Y = (Y_1, \dots, Y_k, Y_{k+1}, \dots, Y_d)$$

Let $Z = (Y_1, \dots, Y_k)$

Generating uniform unit vector:

X_1, \dots, X_d are i.i.d. normal distributions $N(0, 1)$.

$$X = (X_1, \dots, X_d), \quad Y = \frac{1}{\|X\|} X$$

Y is a uniform random unit vector.

The density:

$$\Pr[(x_1, \dots, x_d)] = \prod_{i=1}^d \frac{1}{\sqrt{2\pi}} e^{-x_i^2/2} = (2\pi)^{-d/2} e^{-\|x\|^2/2}$$

Spherically symmetric!

X_1, \dots, X_d are i.i.d. normal distributions $N(0, 1)$.

$$X = (X_1, \dots, X_d), \quad Z = \frac{1}{\|X\|}(X_1, \dots, X_k)$$

$$\|Z\|^2 = \frac{X_1^2 + \dots + X_k^2}{X_1^2 + \dots + X_d^2}$$

$$\begin{aligned} \|Z\|^2 < \frac{k}{2d} &\iff (k - 2d) \sum_{i=1}^k X_i^2 + k \sum_{i=k+1}^d X_i^2 > 0 \\ \|Z\|^2 > \frac{3k}{2d} &\iff (3k - 2d) \sum_{i=1}^k X_i^2 + 3k \sum_{i=k+1}^d X_i^2 < 0 \end{aligned}$$

Sum of independent random variables! Chernoff?

X_1, \dots, X_d are i.i.d. normal distributions $N(0, 1)$.

Upper bound: $\Pr \left[(k - 2d) \sum_{i=1}^k X_i^2 + k \sum_{i=k+1}^d X_i^2 > 0 \right]$

$$(\text{for } \lambda > 0) = \Pr \left[\exp \left\{ \lambda \left((k - 2d) \sum_{i=1}^k X_i^2 + k \sum_{i=k+1}^d X_i^2 \right) \right\} > 1 \right]$$

$$(\mathbf{Markov}) \leq \mathbb{E} \left[\exp \left\{ \lambda \left((k - 2d) \sum_{i=1}^k X_i^2 + k \sum_{i=k+1}^d X_i^2 \right) \right\} \right]$$

$$(\mathbf{independence}) = \prod_{i=1}^k \mathbb{E} \left[e^{\lambda(k-2d)X_i^2} \right] \cdot \prod_{i=k+1}^d \mathbb{E} \left[e^{\lambda k X_i^2} \right]$$

$$E[e^{sx^2}]$$

$$= \frac{1}{\sqrt{2\pi}} \int e^{-(1-2s)x^2/2} dx$$

$$= \frac{1}{\sqrt{1-2s}} \underbrace{\frac{1}{\sqrt{2\pi}} \int e^{-y^2/2} dy}_{\text{Value} = 1} = \frac{1}{\sqrt{1-2s}}$$

tributions $N(0, 1)$.

$$_i^2 + k \sum_{i=k+1}^d X_i^2 > 0$$

$$\left\{ \sum_{i=1}^k X_i^2 + k \sum_{i=k+1}^d X_i^2 \right\} > 1$$

$$X_i^2 + k \sum_{i=k+1}^d X_i^2 \right) \right\}$$

(independence)

$$= \prod_{i=1}^k \mathbf{E} \left[e^{\lambda(k-2d)X_i^2} \right] \cdot \prod_{i=k+1}^d \mathbf{E} \left[e^{\lambda k X_i^2} \right]$$

$$\mathbf{E} \left[e^{sX_i^2} \right] = \frac{1}{\sqrt{1 - 2s}}.$$

$$= (1 - 2\lambda(k - 2d))^{-\frac{k}{2}} (1 - 2\lambda k)^{-\frac{d-k}{2}}$$

$$X_1,\ldots,X_d \text{ are i.i.d. normal distributions } N(0,1).$$

$$\Pr\left[(k-2d)\sum_{i=1}^k X_i^2 + k\sum_{i=k+1}^d X_i^2 > 0\right]$$

$$\leq \left(1 - 2\lambda(k-2d)\right)^{-\frac{k}{2}} \left(1 - 2\lambda k\right)^{-\frac{d-k}{2}} \leq e^{-k/16}$$

$$\text{minimized when }\lambda=\tfrac{1}{2d-k}$$

$$\Pr\left[(3k-2d)\sum_{i=1}^k X_i^2 + 3k\sum_{i=k+1}^d X_i^2 < 0\right] \leq e^{-k/24}$$

$$Z = \tfrac{1}{\|X\|}(X_1,\dots,X_k) \qquad \text{for some $k = O(\ln n)$}$$

$$\Pr\left[\|Z\|^2 > \frac{3k}{2d} \vee \|Z\|^2 < \frac{k}{2d}\right] = o\big(\frac{1}{n^3}\big)$$

for some $k = O(\ln n)$

Z : fixed k -subspace of random unit vector.

$$\Pr [\|Z\|^2 \text{ distorted from } k/d] = o\left(\frac{1}{n^3}\right)$$

Au : random k -subspace of fixed unit vector.

$$\Pr [\|Au\|^2 \text{ distorted from } k/d] = o\left(\frac{1}{n^3}\right)$$

$\forall u, v$

$$\Pr \left[\left\| \sqrt{d/k}A(u - v) \right\|^2 \text{ distorted from } \|u - v\|^2 \right] = o\left(\frac{1}{n^3}\right)$$

$O(n^2)$ pairs of (u, v) . **Union bound \Rightarrow Johnson-Lindenstrauss**

Johnson-Lindenstrauss Theorem

Johnson-Lindenstrauss Theorem:

For any $0 < \epsilon < 1$, for any set V of n points in \mathbf{R}^d ,
there is a map $\phi : \mathbf{R}^d \rightarrow \mathbf{R}^k$ with $k = O(\ln n)$,
such that $\forall u, v \in V$,

$$(1 - \epsilon) \|u - v\|^2 \leq \|\phi(u) - \phi(v)\|^2 \leq (1 + \epsilon) \|u - v\|^2$$

- $\phi(v) = Av$.
- A is a random projection matrix.

Random Projection

Random $k \times d$ matrix A :

- Projection onto a uniform random subspace.

(Johnson-Lindenstrauss)

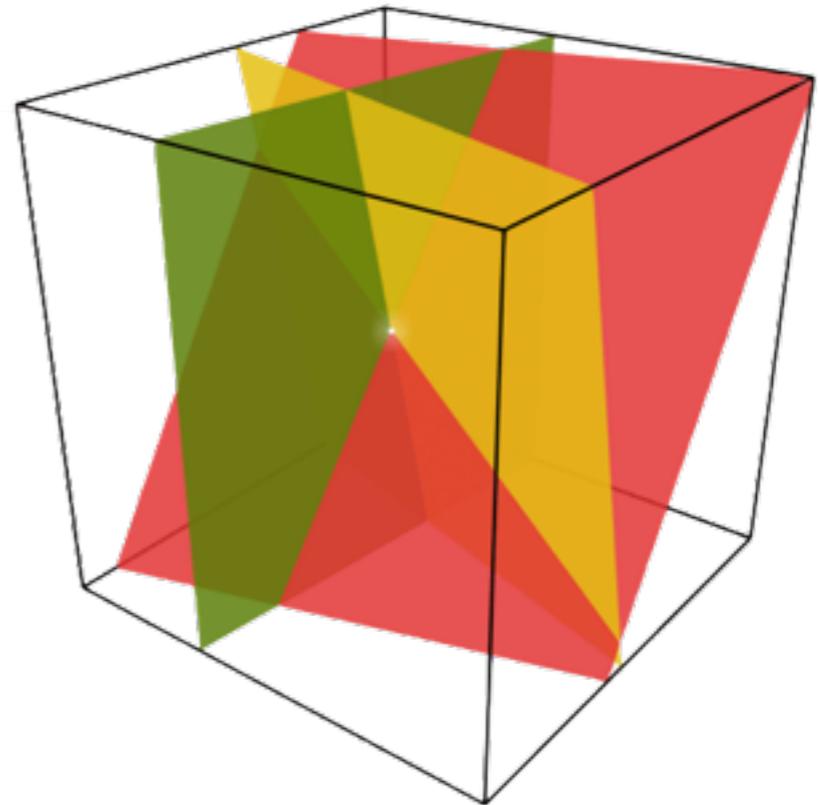
(Dasgupta-Gupta)

- i.i.d. Gaussian entries.

(Indyk-Motiwani)

- i.i.d. -1/+1 entries.

(Achlioptas)



rows: $A_{1\cdot}, A_{2\cdot}, \dots, A_{k\cdot}$.

random orthogonal
unit vectors $\in \mathbb{R}^d$